# On Approximation of 2D Persistence Modules by Interval-decomposables^ 

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#### Abstract

In this work, we propose a new invariant for 2 D persistence modules called the compressed multiplicity and show that it generalizes the notions of the dimension vector and the rank invariant. In addition, for a 2 D persistence module $M$, we propose an "interval-decomposable approximation" $\delta^{*}(M)$ (in the split Grothendieck group of the category of persistence modules), which is expressed by a pair of interval-decomposable modules, that is, its positive and negative parts. We show that $M$ is interval-decomposable if and only if $\delta^{*}(M)$ is equal to $M$ in the split Grothendieck group. Furthermore, even for modules $M$ not necessarily intervaldecomposable, $\delta^{*}(M)$ preserves the dimension vector and the rank invariant of $M$. In addition, we provide an algorithm to compute $\delta^{*}(M)$ (a high-level algorithm in the general case, and a detailed algorithm for the size $2 \times n$ case).


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## 1. Introduction

Persistent homology [1, 2] is one of the main tools in the rapidly growing field of topological data analysis. Given a filtration - a one-parameter increasing sequence of spaces - persistent homology captures the persistence of topological features such as connected components, holes, voids, etc. in the filtration. Here, the persistence of features is quantified by birth and death parameter values. This can be summarized compactly by the so-called persistence diagram, which is the multiset of birth-death pairs drawn on the plane with multiplicity. Algebraically, the persistence diagram can be explained as resulting from a structure theorem (the Krull-Schmidt theorem (Theorem 2.1) and Gabriel's Theorem [3]) of persistence modules, which can also be regarded as representations of certain quivers. See Section 2 for detailed definitions.

One way to deal with multiparametric data is to use multidimensional persistence [4]. However, multidimensional persistence presents theoretical difficulties that hinder the construction of a persistence diagram as in one-dimensional persistence. In particular, there is no complete discrete invariant that captures all isomorphism classes of indecomposable persistence modules [4]. Another way of expressing this difficulty is that the equioriented $m \times n$ commutative grid $\vec{G}_{m, n}$ of sufficiently large size ( $m, n \geq 2$ and $m n \geq 12$, see [5, Theorem 1.3], [6, Theorem 2.5], [7, Theorem 5]) is of wild representation type.

[^0]One way to avoid this problem is to consider only a restricted class of persistence modules. Inspired by 1D persistence, there has been much interest in the so-called interval-decomposable representations, which are direct sums of interval representations (Definition 2.6). The work [8] studied this family of representations and provided a criterion to determine whether or not a given persistence module is interval-decomposable.

It is hoped that most persistence modules coming from "real-world data" contain very few or indeed no non-interval summands. Let us consider the silica glass example computed in [9], which compares the atomic configuration of silica glass with its configuration after physical pressurization. The underlying bound quiver is the commutative ladder $C L_{3}(f b)$, with only two non-interval indecomposable representations given by dimension vectors $\left(\begin{array}{lll}1 & 1 & 1 \\ 1 & 2 & 1\end{array}\right)$ and $\left(\begin{array}{ccc}1 & 2 & 1 \\ 0 & 1 & 0\end{array}\right)$. Then, the numerical result in [9] has $\left(\begin{array}{lll}1 & 1 & 1 \\ 1 & 2 & 1\end{array}\right)$ appearing with only multiplicity 1 and ( $\left.\begin{array}{ccc}1 & 2 & 1 \\ 0 & 1 & 0\end{array}\right)$ with multiplicity 0 , in an example with more than ten thousand indecomposable summands. While in the slightly different setting of a non-equioriented commutative ladder, this provides an example where the non-interval part is minute compared to the interval-decomposable part.

On the other hand, the work [10] argues via a geometric example that the non-interval indecomposables may contain important information that should not be ignored, and that even in relatively simple geometric point clouds embedded in $\mathbb{R}^{3}$, indecomposable summands with arbitrarily large dimension (as a vector space) may be present. These large indecomposable summands are clearly not interval.

In this work, we take neither position, but instead propose a method to replace an arbitrary persistence module $M \in \operatorname{rep} \vec{G}_{m, n}$ by an object $\delta^{*}(M)$ in the split Grothendieck group that is intervaldecomposable. The interval-decomposable approximation $\delta^{*}(M)$ (Definition 5.9) is expressed by a pair of interval-decomposable modules, that is, its positive part $\delta^{*}(M)_{+}$and negative part $\delta^{*}(M)_{-}$(see (5.5)). To construct $\delta^{*}(M)$, we first define what we call the compressed multiplicity (Definition 4.12) of $M$ by a compression operation that picks up information in $M$ restricted to certain essential vertices of intervals.

The intuition behind the compressed multiplicity can be explained as follows. As an initial goal, we want to compute the multiplicity of an interval $I$ as a direct summand of $M$. Indeed, the work [8] presents an algorithm for this computation. However, as this may not be straightforward, in this work we adopt a different approach. We first compress both $M$ and $I$ by restricting the underlying domain to certain essential vertices of $I$, and compute the multiplicity in the representation category with smaller underlying domain.

In the equioriented commutative ladder [9] case $\left(\vec{G}_{2, n}\right)$, the compression operation reduces the underlying bound quiver to a representation-finite bound quiver. This enables easy computation of the compressed multiplicity using preexisting algorithms.

We show that the compressed multiplicity in fact generalizes the notions of dimension vector (Proposition 4.18) and rank invariant (Proposition 4.16). Furthermore, we exhibit representations that can be distinguished by their compressed multiplicities but not by their rank invariants. We thus propose the compressed multiplicity as a new, finer invariant for 2D persistence modules. Moreover, we show that for interval-decomposable representations, the multiplicity can be recovered from the compressed multiplicity (Theorem 4.23).

Then, the object $\delta^{*}(M)$ is defined using the Möbius inversion of the compressed multiplicity of $M$. This is a generalization of the well-known fact that the multiplicities of interval summands in 1D persistence modules can be obtained via an application of inclusion-exclusion on the ranks of the linear maps (see for example $[2,11]$ ).

That is, the persistence diagram is simply the Möbius inversion of the rank invariant. We note that several works have already exploited this observation to define "generalized persistence diagrams" in general settings. In Subsection 1.1, we review some of them and contrast them with our work.

In the case that $M$ is interval-decomposable, it follows that $\delta^{*}(M)$ is equal to $M$ viewed as an element $\llbracket M \rrbracket$ of the split Grothendieck group (Theorem 5.10); that is, $\delta^{*}(M)_{+} \cong M$ and $\delta^{*}(M)_{-}=0$. Furthermore, we show that even for modules $M$ not necessarily interval-decomposable, $\delta^{*}(M)$ preserves the dimension vector and the rank invariant of $M$ (Corollary 5.14, Theorem 5.12). In this sense, we think of $\delta^{*}(M)$ as an interval-decomposable "approximation" of $M$.

We organize this work as follows. In Section 2, we review the necessary background from representation theory and poset theory, and then, in Section 3, we study the poset of interval representations. In Section 4, we introduce our concept of compressed multiplicities and study its properties. In Section 5, we give the construction of $\delta^{*}(M)$ from $M$ via Möbius inversion of the compressed multiplicity and give some results
about its properties. In Section 6, we discuss the computation of our proposed compressed multiplicity and the interval-decomposable approximation.

### 1.1. Möbius inversions in persistence

The work of Patel [12] also uses the idea of Möbius inversion in order to define generalized persistence diagrams, but only in the setting of persistence modules over ( $\mathbb{R}, \leq$ ) [12, Definition 2.1]. In this work, our concept of interval approximation can also seen as an application of Möbius inversion for the more general setting of the 2D commutative grid. However, we do not consider "generalized persistence diagrams" in the sense of [12], but rather restrict our attention to the poset of interval subquivers as we are motivated by their use in practical computation and applications.

On the other hand, the work [13] defines a concept of a "persistence diagram" for $n \mathrm{D}$ persistence modules by using a Möbius inversion in a similar way as we do, and shows a bottleneck stability result. They consider only (hyper)rectangles instead of intervals as the domain for their multiplicity functions and use a partial order specifically tailored for proving bottleneck stability.

After an initial version of this work was sent for review, we were made aware by a reviewer of the prior work of Kim and Memoli [14], which further generalizes Patel's generalized persistence diagram [12]. In Table 1, we provide a rough overview of the different settings and a correspondence of some of the results, which we explain in detail below.

Table 1: Settings and Some Similar Results*

|  | This work | Kim and Memoli [14] |
| :---: | :---: | :---: |
| (1) Underlying setting | commutative grid $\vec{G}_{m, n}$ | poset ${ }^{1} P$ |
| (2) Target category | $\operatorname{vect}_{K}$ | category $^{2} \mathcal{C}$ |
| (3) Domain of invariant | $\mathbb{I}_{m, n}{ }^{3}$ | $\boldsymbol{C o n}(P)^{4}$ |
| (4) Invariant proposed | compressed multiplicities $\underline{d}_{M}^{*}: \mathbb{I}_{m, n} \rightarrow \mathbb{N}$ | generalized rank invariant ${ }^{5}$ <br> $\operatorname{rk}(M): \operatorname{Con}(P) \rightarrow \mathcal{J}(\mathcal{C})$ |
| (5) Inversion | $\delta_{M}^{*}: \mathbb{I}_{m, n} \rightarrow \mathbb{Z}$ | generalized persistence diagram $\operatorname{dgm}^{P}(M): \operatorname{Con}(P) \rightarrow \operatorname{Gr}(\mathcal{C})$ |
| (6) Object | interval-decomposable approximation $\delta^{*}(M) \in \operatorname{Gr}\left(\operatorname{rep} \vec{G}_{m, n}\right)$ | - 6 |
| (7) from proposed invariant to true multiplicities (interval-decomposable) | Theorem 4.23 | [14, Theorem 3.14] |
| from true multiplicities to proposed invariant (interval-decomposable) | Lemma 4.21 | [14, Proposition 3.17] |
| (9) $\begin{aligned} & \text { Interpretation as } \\ & \text { Möbius inversion }\end{aligned}$ | Theorem 5.3 | [14, Proposition 3.19] |

* This table is not intended to be a comprehensive summary of all results. ${ }^{7}$
${ }^{1}$ Essentially finite connected poset
${ }^{2}$ Essentially small, symmetric monoidal category satisfying [14, Convention 2.3]
3 interval (connected and convex) subquivers
${ }^{4}$ path-connected subposets
${ }^{5}$ See [14, Definition 3.5]. The codomain $\mathcal{J}(\mathcal{C})$ is the set of isomorphism classes of $\mathcal{C}$.
${ }^{6}$ Not explicitly defined. See however, [14, Remark 3.22].
${ }^{7}$ For example, [14] contains results concerning Reeb graphs, which can be viewed as functors from the "zigzag poset" to the category of finite sets.

While Kim and Memoli [14] consider a very general setting, we restrict our attention to $K$-representations of the commutative grid $\vec{G}_{m, n}$ (see rows (1) and (2) of Table 1). Since $\vec{G}_{m, n}$ can be viewed as a poset $P$, which happens to be essentially finite and connected, their setting contains ours. However, the domains of
the proposed invariants (see row (3) of Table 1) are different. We note that $\boldsymbol{C o n}(P)$, the set of all pathconnected subposets is in general different from the set of all interval subposets, and this is indeed the case for $P=\vec{G}_{m, n}$. Our use of intervals is motivated by our ultimate goal of constructing an approximation to $M$. In contrast, the set $\operatorname{Con}(P)$ contains subposets which cannot be realized as the support of some persistence module. For example, viewing $\vec{G}_{2,2}$ as a poset with Hasse diagram (both filled and unfilled circles):

the subposet $C$ given by the filled-in circles is in $\boldsymbol{\operatorname { C o n }}(P)$. However, there is no thin ${ }^{1}$ indecomposable persistence module over $\vec{G}_{2,2}$ with support given by $C$, as a commutativity relation will be violated otherwise.

Furthermore, the proposed invariants (row (4) of Table 1) are different. We first note that both papers use of the idea of restricting the input persistence module $M$ to define the respective invariants. In [14], $M$ is restricted to $I \in \mathbf{C o n}(P)$ to obtain $\left.M\right|_{I}$. In the case that $I$ is in fact an interval, this corresponds to applying what we call the "total compression" functor (Definition 4.11) in a more general setting.

Kim and Memoli [14] then defines the value of their generalized rank invariant at $I \in \mathbf{C o n}(P)$ to be "the isomorphism class of the image of the canonical limit-to-colimit map" for $M_{I}$. Of course, in the case that the target category $\mathcal{C}$ is vect $_{K}$, the category of finite-dimensional $K$-vector spaces, this value can be fully characterized by the dimension of the image. In fact, one version of our invariant, which we call the "total compressed multiplicity", coincides with the dimensions of their generalized rank invariant (see Remark 4.13).

Remark 1.1. However, we emphasize that this total compressed multiplicity is not the main emphasis of this work. Instead, motivated by computation, we propose the use of the source-sink (ss-)compression yielding smaller representations (compared to $\left.M\right|_{I}$ ), by further restriction to what we call the essential vertices of $I$. We note that these do not coincide with the generalized rank invariant of [14] for fixed $I$. See Example 4.14. However, if we allow to change the form of the "input" to generalized rank invariant and broaden its domain of definition, we indeed recover our source-sink multiplicity (See Remark 4.15).

## 2. Background

### 2.1. Representation Theory

We first recall some fundamental terminologies of representations of quivers (see [15] for instance ${ }^{2}$ ).
A quiver $Q$ is a quadruple ( $\left.Q_{0}, Q_{1}, s, t\right)$ of sets $Q_{0}, Q_{1}$ of vertices and arrows, respectively and maps $s, t: Q_{1} \rightarrow Q_{0}$ that give the source and target vertices, respectively, of the arrows. We denote an arrow $\alpha$ with source $s(\alpha)=x$ and target $t(\alpha)=y$ by $\alpha: x \rightarrow y$. In this paper, all quivers $Q$ are assumed be finite, namely, $Q_{0}$ and $Q_{1}$ are finite.

Throughout this work, we fix a field $K$. Let $Q$ be a quiver. A representation $V$ of $Q$ (over $K$ ) is a family $(V(x), V(\alpha))$ of a vector space $V(x)$ for each vertex $x \in Q_{0}$ and a linear map $V(\alpha): V(x) \rightarrow V(y)$ for each arrow $\alpha: x \rightarrow y$ in $Q_{1}$.

The dimension vector $\underline{\operatorname{dim}}(V)$ of a representation $V$ of $Q$ is defined as the tuple

$$
\underline{\operatorname{dim}}(V):=(\operatorname{dim} V(x))_{x \in Q_{0}} .
$$

It is customary to display the dimension vector by writing each number $\operatorname{dim} V(x)$ relative to where the vertex $x$ is located on an illustration of the quiver $Q$. The dimension of $V$ is $\operatorname{dim} V:=\sum_{x \in Q_{0}} \operatorname{dim} V(x)$. A

[^1]representation $V$ of $Q$ is said to be finite-dimensional if $\operatorname{dim} V<\infty$. In this work, by representation we mean finite-dimensional representation.

Let $V$ and $W$ be representations of $Q$. A morphism $f: V \rightarrow W$ from $V$ to $W$ is a family $\left(f_{x}\right)_{x \in Q_{0}}$ of linear maps $f_{x}: V(x) \rightarrow W(x)$ such that the following diagram commutes for each arrow $\alpha: x \rightarrow y$ :

$$
\begin{gathered}
V(x) \xrightarrow{f_{x}} W(x) \\
V(\alpha) \downarrow \\
V(y) \xrightarrow[f_{y}]{ } W W(y) .
\end{gathered}
$$

The composition of morphisms $f=\left(f_{x}\right)_{x \in Q_{0}}: V \rightarrow W$ and $g=\left(g_{x}\right)_{x \in Q_{0}}: U \rightarrow V$ is defined in the obvious way: $f \circ g: U \rightarrow W$ is given by $(f \circ g)_{x}=f_{x} \circ g_{x}$. We denote by rep $Q$ the category of finite-dimensional representations of $Q$ together with these morphisms and this composition.

For each vertex $i \in Q_{0}$, we have the path of length 0 at $i$, which is denoted by $e_{i}$. For a given positive integer $n$, a path $\mu$ of length $n$ is a sequence $\alpha_{n} \cdots \alpha_{1}$ of arrows $\alpha_{i}$ such that $t\left(\alpha_{i}\right)=s\left(\alpha_{i+1}\right)$ for all $i=1, \cdots, n-1$. The source vertex of $\mu$ is $s\left(\alpha_{1}\right)$, while its target vertex is $t\left(\alpha_{n}\right)$. An $m$-tuple $\mu_{1}, \cdots, \mu_{m}$ of paths is said to be parallel if they all have the same source vertex and the same target vertex. A relation $\rho$ in $Q$ is a formal sum $\rho=\sum_{i=1}^{m} t_{i} \mu_{i}$ of parallel paths $\mu_{i}$ of length at least 2 with $t_{i} \in K$. A pair $(Q, R)$ of a quiver $Q$ and a set $R$ of relations is called a bound quiver.

A relation $\rho$ is called a commutativity relation if $\rho=\mu_{1}-\mu_{2}$ for some two parallel paths $\mu_{1}, \mu_{2}$. If $R$ is the set of all possible commutative relations in $Q,(Q, R)$ is called a quiver with full commutativity relations.

Let $(Q, R)$ be a bound quiver and let $V$ be a representation of $Q$. Put $V(\mu):=V\left(\alpha_{n}\right) \circ \cdots \circ V\left(\alpha_{1}\right)$ for any path $\mu=\alpha_{n} \cdots \alpha_{1}$ of length $n \geq 1$. Then, $V \in \operatorname{rep} Q$ is said to be a representation of $(Q, R)$ if $V(\rho):=\sum_{i=1}^{m} t_{i} V\left(\mu_{i}\right)=0$ for any $\rho=\sum_{i=1}^{m} t_{i} \mu_{i} \in R$. We denote by $\operatorname{rep}(Q, R)$ the full subcategory of $\operatorname{rep} Q$ consisting of the representations of $(Q, R)$.

The path-category $K Q$ of $Q$ over $K$ is defined as follows. The objects of $K Q$ are the vertices of $Q_{0}$. For each pair $(i, j)$ of objects of $K Q$, the morphisms from $i$ to $j$ are the linear combinations of paths from $i$ to $j$. The composition of $K Q$ is defined as the bilinearization of the concatenation of paths. Then for each object $i$ of $K Q$, the identity morphism of $i$ is given as the path $e_{i}$ of length 0 at $i$. Note that the obtained category $K Q$ naturally becomes a $K$-category, in the sense that $K Q(i, j)$ are $K$-vector spaces for all $i, j \in Q_{0}$, and the composition is $K$-bilinear. The factor category $K Q /\langle R\rangle$ is denoted by $K(Q, R)$, where $\langle R\rangle$ is the ideal of the $K$-category $K Q$ generated by $R$. For instance, this notation is used later for $(Q, R)=\vec{G}_{m, n}$ in Section 4 (see Definition 4.7). For each morphism $\mu$ in $K Q$, the morphism $\mu+\langle R\rangle$ in $K Q /\langle R\rangle$ is usually denoted just by $\mu$, and for morphisms $\mu$ and $\nu$ in $K Q$, we regard $\mu=\nu$ in $K Q /\langle R\rangle$ if and only if $\mu-\nu \in\langle R\rangle$.

A $K$-linear functor from $K(Q, R)$ to vect ${ }_{K}$, the category of finite-dimensional $K$-vector spaces, is called a (left) $K(Q, R)$-module, which can be identified with a representation of $(Q, R)$ in an obvious way. From this fact, representations of $(Q, R)$ are sometimes called modules (over $K(Q, R)$ ).

A fundamental result in representation theory is the Krull-Schmidt theorem (see [16, Theorem 12.9] or [15, I.4.10 Unique decomposition theorem]).

Theorem 2.1 (Krull-Schmidt). Let $\mathcal{L}$ be a complete set of representatives of isomorphism classes of indecomposable representations of a bound quiver $(Q, R)$. For each representation $M$ of $(Q, R)$, there exists a unique function $d_{M}: \mathcal{L} \rightarrow \mathbb{Z}_{\geq 0}$ such that

$$
M \cong \bigoplus_{L \in \mathcal{L}} L^{d_{M}(L)}
$$

The function $d_{M}$ is called the multiplicity function of $M$, and the value $d_{M}(L)$ the multiplicity of the indecomposable $L$ in $M$.

As an example, let us consider the equioriented $A_{n}$-type quiver:

$$
\vec{A}_{n}: 1 \longrightarrow 2 \longrightarrow \cdots \longrightarrow n .
$$

It is known that in this case, $\mathcal{L}$ is the set $\{\mathbb{I}[b, d]\}_{1 \leq b \leq d \leq n}$ of the so-called interval representations $\mathbb{I}[b, d]$ of $\vec{A}_{n}[3]$. The interval representation $\mathbb{I}[b, d]$ is

$$
\mathbb{I}[b, d]: 0 \longrightarrow \cdots \longrightarrow 0 \longrightarrow \stackrel{b-\mathrm{th}}{K} \longrightarrow K \longrightarrow \cdots \longrightarrow \stackrel{d \text {-th }}{K} \longrightarrow 0 \longrightarrow \cdots \longrightarrow 0
$$

which has the vector space $\mathbb{I}[b, d](i)=K$ at the vertices $i$ with $b \leq i \leq d$, and 0 elsewhere, and where the maps between the neighboring vector spaces $K$ are identity maps and zero elsewhere. In the context of persistent homology [1, 2], a persistence module can be viewed as a representation of $\vec{A}_{n}$, and the multiplicity function $d_{M}$ encodes the information of the persistence diagram.

The underlying bound quiver we study in this work is the equioriented commutative grid $\vec{G}_{m, n}$ defined below. Then, we consider 2D persistence modules as representations of $\vec{G}_{m, n}$.

Definition 2.2 (Equioriented commutative grid). Let $0<m, n \in \mathbb{Z}$. The bound quiver $\vec{G}_{m, n}$, is defined to be the 2D grid of size $m \times n$ with all horizontal arrows in the same direction and all vertical arrows in the same direction, together with full commutativity relations. It is also called the equioriented commutative grid of size $m \times n$.

For example, the equioriented $2 \times 4$ commutative grid $\vec{G}_{2,4}$ is the quiver

with full commutativity relations.
As mentioned in the introduction, for large enough size, $\vec{G}_{m, n}$ is of wild representation type. That is, $\mathcal{L}$ can be very complicated. Instead, we consider a restricted class of representations, the interval-decomposable representations. Following the notation in [8], we first recall the definition of interval subquivers and interval representations for general bound quivers.
Definition 2.3 (Interval subquiver).
(1) Let $Q$ be a quiver. A full subquiver $Q^{\prime}$ of $Q$ is said to be convex in $Q$ if and only if for all vertices $x$, $y \in Q_{0}^{\prime}$ and all vertices $z \in Q_{0}$, the existence of paths $x$ to $z$ and $z$ to $y$ in $Q$ imply that $z \in Q_{0}^{\prime}$.
(2) A quiver Q is said to be connected if it is connected as an "undirected graph",
(3) A subquiver $Q^{\prime}$ of $Q$ is said to be an interval subquiver of $Q$ if $Q^{\prime}$ is convex (in $Q$ ) and connected.

For any two full subquivers $Q^{\prime}, Q^{\prime \prime}$ of $Q$, the intersection $Q^{\prime} \cap Q^{\prime \prime}$ (respectively, the union $Q^{\prime} \cup Q^{\prime \prime}$ ) of $Q^{\prime}$ and $Q^{\prime \prime}$ is defined as the full subquiver of $Q$ having the vertex set $Q_{0}^{\prime} \cap Q_{0}^{\prime \prime}$ (respectively, $Q_{0}^{\prime} \cup Q_{0}^{\prime \prime}$ ).

Suppose that $Q^{\prime}$ and $Q^{\prime \prime}$ are interval subquivers of $Q$ with $Q_{0}^{\prime} \cap Q_{0}^{\prime \prime} \neq \emptyset$. Note that $Q^{\prime} \cap Q^{\prime \prime}$ may not be connected, in general, and so may not be an interval. However, the following statement can be checked.

Lemma 2.4. Let $Q^{\prime}$ and $Q^{\prime \prime}$ be interval subquivers of $Q$. Then, $Q^{\prime} \cap Q^{\prime \prime}$ is a disjoint union of interval subquivers of $Q$.

Proof. To see this, we write $Q^{\prime} \cap Q^{\prime \prime}$ as a disjoint union of its connected components $C_{i}$ for $i=1, \cdots, n$ and show that each connected component $C_{i}$ is actually an interval subquiver of $Q$. It suffices to check that $C_{i}$ is convex.

For that, let $x, y$ be vertices of $C_{i}$ and $z$ a vertex of $Q$ such that there exist paths $x$ to $z$ and $z$ to $y$ in $Q$. We show that $z$ is a vertex of $C_{i}$.

For each path $z=z_{0} \rightarrow z_{1} \rightarrow \cdots \rightarrow z_{\ell}=y$ in $Q$, since $Q^{\prime}$ and $Q^{\prime \prime}$ are convex and $x, y$ are both in $Q^{\prime}$ and $Q^{\prime \prime}$, each $z_{k}$ is a vertex of $Q^{\prime}$ and $Q^{\prime \prime}$. Thus, all $z_{k}$ are vertices in $Q^{\prime} \cap Q^{\prime \prime}$ and the path $z=z_{0} \rightarrow z_{1} \rightarrow \cdots \rightarrow z_{\ell}=y$ is in fact a path in $Q^{\prime} \cap Q^{\prime \prime}$. Since $z_{\ell}=y \in C_{i}$ and $C_{i}$ is a connected component, we must have $z_{0}=z \in C_{i}$. Thus $C_{i}$ is convex.

On the other hand, $Q^{\prime} \cup Q^{\prime \prime}$ is not an interval subquiver in general, even if $Q^{\prime}$ and $Q^{\prime \prime}$ are interval subquivers of $Q$ with $Q_{0}^{\prime} \cap Q_{0}^{\prime \prime} \neq \emptyset$. While connectedness is guaranteed since $Q_{0}^{\prime} \cap Q_{0}^{\prime \prime} \neq \emptyset$, convexity may fail to hold.

Definition 2.5. For $0<m, n \in \mathbb{Z}$, define $\mathbb{I}_{m, n}$ to be the set of all nonempty interval subquivers of $\vec{G}_{m, n}$.
It is known that the interval subquivers of $\vec{G}_{m, n}$ take on a distinctive "staircase" shape. See [8]. Below is an example of an interval subquiver of $\vec{G}_{4,6}$.


Recall that for $M$ a representation of a bound quiver $(Q, R)$, the support $\operatorname{supp} M$ of $M$ is the full subquiver of $Q$ with vertices $\{i \in Q \mid M(i) \neq 0\}$. Finally, we are ready to recall the following generalization of interval representations of $\vec{A}_{n}$.

Definition 2.6 (Interval representations). A representation $V \in \operatorname{rep}(Q, R)$ is said to be an interval representation if

- $\operatorname{dim} V(x) \leq 1$ for each vertex $x$ of $Q$,
- its $\operatorname{support} \operatorname{supp}(V)$ is an interval of $Q$, and
- for all arrows $\alpha \in \operatorname{supp}(V), V(\alpha)$ is an identity map.

Note that by this definition, an interval representation $V$ is determined (up to isomorphism) by its support supp $V$. If $I$ is an interval subquiver, the corresponding interval representation with support equal to $I$ is denoted by $V_{I}$. For example, the interval subquiver $I$ of $\vec{G}_{4,6}$ given by the quiver (2.1) is the support of $V_{I}$ with dimension vector (3.2).

A representation $M \in \operatorname{rep}(Q, R)$ is said to be interval-decomposable if it can be expressed as a direct sum of interval representations. Equivalently, by Theorem 2.1, $M$ is interval-decomposable if and only if $d_{M}(L)=0$ for all non-interval indecomposables $L$.

### 2.2. Posets and Lattices

In this subsection, we recall some basic definitions from poset and lattice theory. See [17] for more details.
Recall that a poset (partially ordered set) $(P, \leq)$ is a set $P$ with partial order $\leq$. A poset $P$ is said to be finite if $P$ is finite as a set. Throughout this work, all posets are assumed to be finite.

Definition 2.7. Let $P$ be a poset and $x, y \in P$. The segment $[x, y]$ between $x$ and $y$ is defined to be

$$
[x, y]:=\{z \in P \mid x \leq z \leq y\}
$$

and define $\operatorname{Seg}(P)$ to be the set of all segments of $P$. The open segment $(x, y)$ between $x$ and $y$ is defined to be

$$
(x, y):=\{z \in P \mid x<z<y\}
$$

It is clear that each segment of $P$ (respectively each open segment) of $P$ forms a subposet of $P$. We say that $y$ covers $x$ if $x<y$ and $(x, y)=\emptyset$. The set of the elements covering $x$ is denoted by $\operatorname{Cov}(x)$.

We note that a segment $[x, y]$ is also called an interval in the literature, but we do not use this term to avoid confusion.

Definition 2.8. Let $P$ be a poset and $S$ a subset of $P$.
(1) An element $u \in P$ is said to be an upper bound of $S$ if $s \leq u$ for each $s \in S$. The set of upper bounds of $S$ is denoted by $U(S)$. For a singleton $S=\{s\}$, we abuse the notation and write $U(s)$ for $U(\{s\})$.
(2) An element $x \in U(S)$ is said to be the join of $S$ if $x \leq u$ for each $u \in U(S)$. Note that the join of $S$ is unique if it exists, and is denoted by $\bigvee S$. When $S=\{a, b\}$, then the join of $S$ is denoted by $a \vee b$.

Dually,
(3) An element $l \in P$ is said to be an lower bound of $S$ if $l \leq s$ for each $s \in S$. The set of lower bounds of $S$ is denoted by $L(S)$. For a singleton $S=\{s\}$, we abuse the notation and write $L(s)$ for $L(\{s\})$.
(4) An element $x \in L(S)$ is said to be the meet of $S$ if $l \leq x$ for each $l \in L(S)$. Note that the meet of $S$ is unique if it exists, and is denoted by $\bigwedge S$. When $S=\{a, b\}$, then the meet of $S$ is denoted by $a \wedge b$.

Definition 2.9. Let $P$ be a poset.
(1) $P$ is called a join-semilattice (respectively, meet-semilattice) if each two-element subset $\{a, b\} \subseteq P$ has a join (respectively, meet).
(2) $P$ is called a lattice if $P$ is a join-semilattice and a meet-semilattice.
(3) When $P$ is a lattice, $P$ is said to be distributive if for all $x, y, z \in P$,

$$
x \wedge(y \vee z)=(x \wedge y) \vee(x \wedge z)
$$

or equivalently, if for all $x, y, z \in P$,

$$
x \vee(y \wedge z)=(x \vee y) \wedge(x \vee z)
$$

For a join-semilattice $P$ and $a, b, c \in P$, note that $(a \vee b) \vee c=\bigvee\{a, b, c\}=a \vee(b \vee c)$. Thus the binary operation $\vee$ satisfies associativity, and hence generalized associativity. Therefore in general, if $S=$ $\left\{x_{1}, \ldots, x_{n}\right\} \subset P$, then

$$
x_{1} \vee x_{2} \vee \cdots \vee x_{n}
$$

is well-defined and equal to $\bigvee S$. A similar remark holds for $\Lambda S$ in meet-semilattices.
The following fact is well-known and can be checked easily.
Proposition 2.10. If $P$ is a finite join-semilattice (meet-semilattice) with a lower bound (upper bound) of $P$, then $P$ is a lattice.

We will see later that the poset of intervals does not form a lattice globally, so we provide the following "local" definitions.

## Definition 2.11.

(1) A poset $P$ is called a local lattice if for any $x, y \in P$, the segment $[x, y]$ is a lattice.
(2) A local lattice $P$ is said to be locally distributive if for any $x, y \in P$, the segment $[x, y]$ is a distributive lattice.

### 2.3. Möbius Functions

In this subsection, we review some basic facts about Möbius functions. We refer the reader again to [17] for more details.

Let $F$ be a field of characteristic zero, and $P$ a poset. Recall that $\operatorname{Seg}(P)$ is the set of segments of $P$. The incidence algebra of $P$ over $F$ is the set of functions from $\operatorname{Seg}(P)$ to $F$, together with a "pointwise" +
operation, and convolution $*$ as the multiplication operation. More precisely, for $f, g: \operatorname{Seg}(P) \rightarrow F$, define $f * g: \operatorname{Seg}(P) \rightarrow F$ by

$$
(f * g)([x, y]):=\sum_{x \leq z \leq y} f([x, z]) g([z, y])
$$

It can be shown that the incidence algebra of $P$ over $F$ is indeed a $F$-algebra, which we denote by $I(P)$. Its identity element is the delta function $\delta: \operatorname{Seg}(P) \rightarrow F$ with

$$
\delta([x, y])= \begin{cases}1_{F} & \text { if } x=y \\ 0 & \text { otherwise }\end{cases}
$$

Definition 2.12 (Zeta and Möbius functions). The zeta function $\zeta: \operatorname{Seg}(P) \rightarrow F$ is the function with constant value $1_{F}$. Then, it can be shown that $\zeta$ is an invertible element of $I(P)$, with inverse called the Möbius function $\mu$.

Now, let $F^{P}$ be the set of all functions $P \rightarrow F$. Note that $F^{P}$ has natural $F$-vector space structure by pointwise addition and scalar multiplication of functions. The incidence algebra $I(P)$ acts on $F^{P}$ from the left by the following. For each $f \in F^{P}, \phi \in I(P)$, define $\phi f \in F^{P}$ by

$$
(\phi f)(x):=\sum_{x \leq y} \phi([x, y]) f(y)
$$

It can be checked that $F^{P}$ is a left $I(P)$ module with this left action. For example, the computation

$$
\begin{aligned}
(\psi(\phi f))(x) & =\sum_{x \leq y} \psi([x, y])(\phi f)(y) \\
& =\sum_{x \leq y} \psi([x, y]) \sum_{y \leq z} \phi([y, z]) f(z) \\
& =\sum_{x \leq z} \sum_{x \leq y \leq z} \psi([x, y]) \phi([y, z]) f(z) \\
& =\sum_{x \leq z}(\psi * \phi)([x, z]) f(z) \\
& =[(\psi * \phi) f](x)
\end{aligned}
$$

shows that this action is compatible with the multiplication (convolution) in $I(P)$.

## 3. Local lattice of intervals

In this section, we study the set of isomorphism classes of interval representations for a fixed equioriented commutative 2D grid $\vec{G}_{m, n}$. Note that an interval representation is uniquely defined (up to isomorphism) by its support, and thus it suffices to consider the set of interval subquivers $\mathbb{I}_{m, n}$.

First, we start with the following easy observation.
Proposition 3.1. With the order $\leq$ on $\mathbb{I}_{m, n}$ defined by $I \leq I^{\prime} \Longleftrightarrow I \subseteq I^{\prime},\left(\mathbb{I}_{m, n}, \leq\right)$ is a poset.
Proof. This is immediate from the definitions.
By Proposition 4.1 in [8], each element $I$ of $\mathbb{I}_{m, n}$ has a "staircase" form, which was denoted by:

$$
I=\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i}
$$

for some integers $1 \leq s \leq t \leq m$ and some integers $1 \leq b_{i} \leq d_{i} \leq n$ for each $s \leq i \leq t$ such that

$$
\begin{equation*}
b_{i+1} \leq b_{i} \leq d_{i+1} \leq d_{i} \text { for all } i \in\{s, \ldots, t-1\} \tag{3.1}
\end{equation*}
$$

In this notation, each $\left[b_{i}, d_{i}\right]_{i}$ is the "slice" of the staircase at height $i$. For example, the staircase

$$
\left(\begin{array}{lllllll}
0 & 1 & 1 & 1 & 0 & 0  \tag{3.2}\\
0 & 0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 1 & 1 \\
0
\end{array}\right)
$$

of $\vec{G}_{4,6}$ corresponds to $[5,6]_{1} \sqcup[3,5]_{2} \sqcup[3,4]_{3} \sqcup[2,4]_{4}$. In general, the interval $I=\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i}$ means that $I$ has vertices

$$
I_{0}=\left\{(i, x) \mid s \leq i \leq t, b_{i} \leq x \leq d_{i}\right\}
$$

Proposition 3.2. Let $I \in \mathbb{I}_{m, n}$ and $J \in \operatorname{Cov}(I)$. Then, the number of vertices of $J$ is one more than that of $I$.

Sketch of Proof. Suppose that $I \subsetneq J$. We show that there exists a point $p \in J_{0} \backslash I_{0}$ that can be added to $I$ to obtain an interval $I^{\prime}$ with $I \subsetneq I^{\prime} \subseteq J$. The result immediately follows from this, since if $J \in \operatorname{Cov}(I)$, then $J=I^{\prime}$ by definition. That is, $J$ has one more vertex compared to $I$.

Let

$$
I=\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i} \text { and } J=\bigsqcup_{j=u}^{v}\left[c_{j}, e_{j}\right]_{j}
$$

Since $I \subsetneq J$, it follows that $u \leq s \leq t \leq v$ and $c_{k} \leq b_{k} \leq d_{k} \leq e_{k}$ for each $k \in[s, t]$, in addition to the requirements for $I$ and $J$ to be intervals. We give below the point $p \in J_{0} \backslash I_{0}$ that can be added to $I$ to obtain the interval $I^{\prime}$.

- In case that $1 \leq u<s$,
- if $c_{s-1} \leq d_{s}$, then choose the point $p=\left(s-1, d_{s}\right)$;
- otherwise, if $c_{s-1}>d_{s}$, choose $p=\left(s, d_{s}+1\right)$.
- The case $t<v \leq m$ is dual to the previous case.
- If $b_{t} \leq e_{t+1}$ choose $p=\left(t+1, b_{t}\right)$;
- otherwise, $p=\left(t, b_{t}-1\right)$ works.
- Otherwise, we have $u=s \leq t=v$. In this case, we define

$$
L=\left\{k \in[s, t] \mid\left(k, b_{k}-1\right) \in J_{0}\right\} \text { and } R=\left\{k \in[s, t] \mid\left(k, d_{k}+1\right) \in J_{0}\right\}
$$

These are the row indices where a point to the left (and right, respectively) of $I$ is in $J$. Since $I \neq J$, it is clear that at least one of $L$ and $R$ is nonempty.

- If $L \neq \emptyset$, choose the point $p=\left(\max L, b_{\max } L-1\right)$.
- If $R \neq \emptyset$, choose the point $p=\left(\min R, d_{\min } R+1\right)$.

For each of the cases above (which exhausts all possibilities), a routine check using the definitions shows that the chosen point $p$ can be added to $I$ to obtain an interval $I^{\prime}$. This completes the proof.

The above result implies that $\mathbb{I}_{m, n}$ is a graded poset with rank function $\rho: \mathbb{I}_{m, n} \rightarrow \mathbb{N}$ given by $\rho(I)=\# I_{0}$, the number of vertices of $I$.

Example 3.3. For any $n \in \mathbb{N}$ and any interval $I=\left[b_{1}, d_{1}\right]_{1} \sqcup\left[b_{2}, d_{2}\right]_{2} \in \mathbb{I}_{2, n}, \# \operatorname{Cov}(I) \leq 4$. Indeed, any cover of $I$ takes on one of the following forms:

$$
\begin{aligned}
& {\left[b_{1}-1, d_{1}\right]_{1} \sqcup\left[b_{2}, d_{2}\right]_{2},} \\
& {\left[b_{1}, d_{1}+1\right]_{1} \sqcup\left[b_{2}, d_{2}\right]_{2},} \\
& {\left[b_{1}, d_{1}\right]_{1} \sqcup\left[b_{2}-1, d_{2}\right]_{2}, \text { or }} \\
& {\left[b_{1}, d_{1}\right]_{1} \sqcup\left[b_{2}, d_{2}+1\right]_{2} .}
\end{aligned}
$$

In general, we have the following, which follows immediately from Proposition 3.2 and the characterization of interval subquivers of $\vec{G}_{m . n}$ as staircases.
Proposition 3.4. Let $I \in \mathbb{I}_{m, n}$. Then, $\operatorname{Cov}(I)=C \cap \mathbb{I}_{m, n}$ where $C$ is the set of subquivers of $\vec{G}_{m, n}$ obtained from I by one of the following operations (if the result is a subquiver):
(1) extending one row of I by one adjacent vertex left of the row,
(2) extending one row of I by one adjacent vertex right of the row,
(3) adding one vertex above the upper-left vertex of $I$, or
(4) adding one vertex below the lower-right vertex of $I$.

Let us express the above using the notation of

$$
I=\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i}
$$

for some integers $1 \leq s \leq t \leq m$ and some integers $1 \leq b_{i} \leq d_{i} \leq n$ for each $s \leq i \leq t$ such that $b_{i+1} \leq b_{i} \leq d_{i+1} \leq d_{i}$ for any $i \in\{s, \ldots, t-1\}$. Then $\operatorname{Cov}(I)$ is the set of valid interval subquivers in the following set of candidates $C$ :

- for $j \in\{s, \ldots, t\}$,

$$
\bigsqcup_{i=s}^{t}\left[b_{i}^{\prime}, d_{i}\right]_{i}, \text { where } b_{i}^{\prime}= \begin{cases}b_{i}-1 & \text { if } i=j \\ b_{i} & \text { otherwise }\end{cases}
$$

- for $j \in\{s, \ldots, t\}$,

$$
\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}^{\prime}\right]_{i}, \text { where } d_{i}^{\prime}= \begin{cases}d_{i}+1 & \text { if } i=j \\ d_{i} & \text { otherwise }\end{cases}
$$

- $\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i} \sqcup\left[b_{t}, b_{t}\right]_{t+1}$,
- $\left[d_{s}, d_{s}\right]_{s-1} \sqcup \bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i}$.

Note that some candidates may exceed the bounds of the commutative grid. Those candidates are immediately disqualified.

Example 3.5. We provide an example using the interval $I$ (filled-in circles):

in the commutative grid $\vec{G}_{5,6}$. We illustrate the vertices in Proposition 3.4.

- Vertices $v$ with $I_{0} \cup\{v\}=C_{0}$ for some $C \in \operatorname{Cov} I$ are denoted with green check marks. These give all the cover elements $C$.
- The remaining vertices $v$ do not form cover elements. That is, there is no interval $C$ with $C_{0}=I_{0} \cup\{v\}$. These are denoted with red crosses. Note that two of them go out of bounds.

$x$
Repeating the point above, each $C \in \operatorname{Cov}(I)$ is the unique interval subquiver $C$ with $I_{0} \cup\{v\}=C_{0}$ for some vertex $v$ given by the green check marks.

Proposition 3.6. The poset $\mathbb{I}_{m, n}$ is a local lattice.
Proof. Let $I, J$ be an intervals of $\mathbb{I}_{m, n}$ with $I \leq J$. We show that the segment $[I, J]$ is a lattice.
Let $J_{1}, J_{2} \in[I, J]$. Then, by Lemma 2.4, the intersection $J_{1} \cap J_{2}$ is given by the disjoint union of some intervals $C_{i}: \bigsqcup_{i=1}^{l} C_{i}$. In this setting, there exists a unique $j$ such that $C_{j}$ contains $I$. Then the meet $J_{1} \wedge J_{2}$ of $J_{1}$ and $J_{2}$ in the segment $[I, J]$ is exactly the interval $C_{j}$. Proposition 2.10 implies that the segment $[I, J]$ is a lattice.

Note that in the above argument, the interval $J$ did not play any role in determining the meet in $[I, J]$. We could have replaced $J$ by the maximum element $M$ in $\mathbb{I}_{m, n}$, which is the entire quiver of $\vec{G}_{m, n}$. That is, the meet of $J_{1}, J_{2}$ in $[I, J]$ is the same as the meet of $J_{1}, J_{2}$ in $[I, M]=U(I)$. Thus, we also call the meet of $J_{1}, J_{2}$ in $[I, J]$ as the meet of $J_{1}, J_{2}$ over $I$.

On the other hand, the join $J_{1} \vee J_{2}$ in $[I, J]$ is the minimum interval containing $J_{1} \cup J_{2}$ by definition. Clearly, $J_{1} \cup J_{2} \subset J \subset M$, and so the join of $J_{1}, J_{2}$ in $[I, J]$ is the same as the join of $J_{1}, J_{2}$ in $[I, M]=U(I)$. Thus, we also call the join of $J_{1}, J_{2}$ in $[I, J]$ as the join of $J_{1}, J_{2}$ over $I$.

Example 3.7. Let $I=\left(\begin{array}{lll}0 & 1 & 0 \\ 0 & 0 & 0\end{array}\right)$ be an interval of $\mathbb{I}_{2,3}$. The intervals $J=\left(\begin{array}{lll}0 & 1 & 1 \\ 0 & 0 & 0\end{array}\right), J^{\prime}=\left(\begin{array}{lll}0 & 1 & 0 \\ 0 & 1 & 0\end{array}\right)$ in $U(I)$ have join $J \vee J^{\prime}=\left(\begin{array}{lll}0 & 1 & 1 \\ 0 & 1 & 1\end{array}\right)$ over $I$.

While we have seen in Proposition 3.6 that $\mathbb{I}_{m, n}$ is a local lattice, it is not a lattice as a whole (Example 3.8), nor is it locally distributive (Example 3.9).

Example 3.8. In general, the meet and join is ill-defined. For example, let $J=\left(\begin{array}{lll}1 & 0 & 0 \\ 0 & 0 & 0\end{array}\right)$ and $J^{\prime}=\left(\begin{array}{lll}0 & 0 & 0 \\ 0 & 0 & 1\end{array}\right)$ be intervals in $\mathbb{I}_{2,3}$. We note that $J \cap J^{\prime}=\emptyset$, so that there is no $I \in \mathbb{I}_{m, n}$ with $J, J^{\prime} \in U(I)$. Then, $X_{1}=\left(\begin{array}{lll}1 & 1 & 1 \\ 0 & 1 & 1\end{array}\right)$ and $X_{2}=\left(\begin{array}{lll}1 & 0 & 0 \\ 1 & 1 & 1\end{array}\right)$ are both minimal among intervals containing both $J$ and $J^{\prime}$. Thus, $J \vee J^{\prime}$, which is supposed to be the minimum interval containing $J \cup J^{\prime}$, is not well-defined. The poset $\mathbb{I}_{m, n}$ is not a lattice, in general.

Example 3.9. In general, the local lattice $\mathbb{I}_{m, n}$ is not locally distributive. Indeed, let $I=\left(\begin{array}{llll}0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0\end{array}\right)$ and $J=\left(\begin{array}{lll}1 & 1 & 1 \\ 1 & 1 & 1\end{array}\right)$ be intervals of $\mathbb{I}_{2,4}$. Moreover, let $I_{1}=\left(\begin{array}{llll}1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0\end{array}\right), I_{2}=\left(\begin{array}{lll}0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 1\end{array}\right)$, and $I_{3}=\left(\begin{array}{lll}0 & 1 & 1 \\ 0 & 0 & 1\end{array}\right)$ be intervals of the segment $[I, J]$. Then we compute $I_{1} \vee\left(I_{2} \wedge I_{3}\right)=I_{1}$ and $\left(I_{1} \vee I_{2}\right) \wedge\left(I_{1} \vee I_{3}\right)=\left(\begin{array}{lll}1 & 1 & 0 \\ 1 & 1 & 1 \\ 0\end{array}\right) \neq I_{1}$.

## 4. Compression and Compressed Multiplicities

In this section, we present the underlying mechanism for an "interval-approximation" that we define and study in Section 5. Here, we define compression functors based on certain essential vertices. These
compression functors then lead to what we call compressed multiplicities. We show that the well-known dimension vector and rank invariant are in fact special cases of compressed multiplicities. Furthermore, we show that for interval-decomposable representations, the true multiplicity information can be recovered from the compressed multiplicies.

### 4.1. Essential Vertices

First, we define two types of "essential vertices".
Recall that a vertex $x$ is said to be a source if there are no arrows $\alpha$ with target $t(\alpha)=x$, and is said to be a $\operatorname{sink}$ if there are no arrows $\alpha$ with source $s(\alpha)=x$.

Definition 4.1 (Source-sink-essential vertices). Let $I$ be an interval subquiver of $\vec{G}_{m, n}$. A vertex $x \in I_{0}$ is said to be source-sink-essential (ss-essential) if $x$ is a source or a sink in $I$.

The set of ss-essential vertices of $I$ will be denoted by $I_{0}^{\text {ss }}$.
Example 4.2. In the following interval subquiver $I$ in $\vec{G}_{6,4}$ :

the vertices denoted by $\circledast$ are ss-essential vertices of $I$.
Lemma 4.3. Let $I, J$ be intervals of $\mathbb{I}_{m, n}$. Assume that $I_{0}^{s s} \subseteq J_{0}$. Then we have $I \leq J$.
Proof. Let $x \in I_{0}$. Then, there is a source $y$, a sink $z$, and a path $\mu$ in $I$ from $y$ to $z$ such that $\mu$ passes through $x$. Since $y, z \in I_{0}^{\text {ss }} \subseteq J_{0}$ and $J$ is convex, we have $x \in J_{0}$, as desired.

Definition 4.4 (Corner-complete-essential vertices). Let $I$ be an interval subquiver of $\vec{G}_{m, n}$. A vertex $x \in I_{0}$ is said to be corner-complete-essential (cc-essential) if $x \in\left(\operatorname{pr}_{1} I_{0}^{\mathrm{ss}} \times \mathrm{pr}_{2} I_{0}^{\mathrm{ss}}\right) \cap I_{0}$, where $\mathrm{pr}_{i}: \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$ is the projection map to the $i$-th axis.

The set of cc-essential vertices of $I$ is denoted by $I_{0}^{\mathrm{cc}}$.
Example 4.5. For the interval subquiver $I$ used in Example 4.2:

the vertices denoted by $\circledast$ are cc-essential vertices of $I$.
Lemma 4.6. Let $I, J$ be intervals of $\mathbb{I}_{m, n}$. Assume that $I_{0}^{c c} \subseteq J_{0}$. Then we have $I \leq J$.
Proof. Since $I_{0}^{\mathrm{ss}} \subseteq I_{0}^{\mathrm{cc}} \subseteq J_{0}$, we have $I \leq J$ by Lemma 4.3.

### 4.2. Compression

In this subsection, we treat both types of essential vertices in parallel to define two types of compression of representations of the equioriented 2D commutative grid $\vec{G}_{m, n}=(Q, R)$. In the previous subsection, we defined the sets of essential vertices $I_{0}^{\mathrm{ss}}$ and $I_{0}^{\mathrm{cc}}$. We consider the full subcategories of $K \vec{G}_{m, n}=K(Q, R)=$ $K Q /\langle R\rangle$ they induce.

Definition 4.7 (ss-compressed category and cc-compressed category). Let $I$ be an interval subquiver of $\vec{G}_{m, n}$ and $E$ be the set of all ss-essential vertices (or cc-essential vertices, respectively) of $I$. The ss-compressed category $I^{\text {ss }}$ (resp. cc-compressed category $I^{\text {cc }}$ ) of $I$ is the full subcategory of $K \vec{G}_{m, n}$ with set of objects $E$.

For completeness, we also introduce the following concept, where we take all vertices of $I$ to be essential. We use the designation "tot" to stand for "total", since all vertices are considered essential in $I^{\text {tot }}$.
Definition 4.8 (compressed category). The compressed category $I^{\text {tot }}$ is the full subcategory of $K \vec{G}_{m, n}$ consisting of all vertices of $I$.

Remark 4.9. For an interval $I$, we distinguish the following similar but different notions related to $I$ : $I$ itself as a full subquiver of $\vec{G}_{m, n}, V_{I}$ the representation of $K \vec{G}_{m, n}$ with support $I$, and $I^{\text {tot }}$ as the full subcategory of $K \vec{G}_{m, n}$ with objects the vertices of $I$.

We note that the bound quiver of $I^{\text {tot }}$ is $\left(I, R_{I}\right)$ with the set of full commutativity relations $R_{I}$. The ss-compressed category or cc-compressed category can also be expressed as a bound quiver, and we identify $\operatorname{rep}\left(Q_{I}^{*}, R_{I}^{*}\right) \cong \operatorname{rep} I^{*}$, where $\left(Q_{I}^{*}, R_{I}^{*}\right)$ is the bound quiver of the compressed category $I^{*}$ for $*=\mathrm{ss}$, cc, tot.

Throughout the rest of this work, we shall use the symbol ' $*$ ' to stand for either 'ss', 'cc' or 'tot' for statements that apply to all three cases as long as it does not cause any confusion.

Example 4.10. For the interval subquiver $I$ in Example 4.2, the compressed categories (displayed as bound quivers) are the following:

and

while


Definition 4.11 (Compression functor). Let $I$ be an interval subquiver of $\vec{G}_{m, n}$ and let $\iota_{I}^{\mathrm{ss}}: I^{\mathrm{ss}} \hookrightarrow K \vec{G}_{m, n}$ (or $\iota_{I}^{\text {cc }}: I^{\text {cc }} \hookrightarrow K \vec{G}_{m, n}$, or $\iota_{I}^{\text {tot }}: I^{\text {tot }} \hookrightarrow K \vec{G}_{m, n}$, respectively) be the inclusion functor into the equioriented 2D commutative grid.

The ss-compression functor $\operatorname{Comp}_{I}^{\mathrm{ss}}(-): \operatorname{rep} K \vec{G}_{m, n} \rightarrow \operatorname{rep} I^{\mathrm{ss}}$ (the cc-compression functor $\operatorname{Comp}_{I}^{\mathrm{cc}}(-)$ or the tot-compression functor $\operatorname{Comp}_{I}^{\mathrm{tot}}(-)$, respectively) is defined by $\operatorname{Comp}_{I}^{\mathrm{ss}}(M)=M \circ \iota_{I}^{\mathrm{ss}}\left(\operatorname{Comp}_{I}^{\mathrm{cc}}(M)=\right.$ $M \circ \iota_{I}^{\text {cc }}$ or $\operatorname{Comp}_{I}^{\text {tot }}(M)=M \circ \iota_{I}^{\text {tot }}$, respectively).

Note that these functors are exactly the restriction functors.
It is clear that the ss-compression, cc-compression and tot-compression functors are additive by definition. To simplify the notation, we let $\operatorname{Comp}_{I}^{*}(-)$ stand for $\operatorname{Comp}_{I}^{\mathrm{ss}}(-), \operatorname{Comp}_{I}^{\mathrm{cc}}(-)$, or $\operatorname{Comp}_{I}^{\mathrm{tot}}(-)$ for statements that hold for all three versions of compression.

Given $M \in \operatorname{rep} \vec{G}_{m, n}$, the compressed representation $\operatorname{Comp}_{I}^{*}(M)$ is a representation of $I^{*}$. Similary, the interval representation $V_{I}$ associated to the interval $I$ has a compressed representation $\operatorname{Comp}_{I}^{*}\left(V_{I}\right)$. For example, the interval $I$ in Example 4.2 is associated to the interval representation

which has ss-compressed representation

a representation of $I^{\mathrm{ss}}$.
While the compressed representation $\operatorname{Comp}_{I}^{*}(M)$ may be interesting in its own right, in the next definition we only consider the multiplicity of $\operatorname{Comp}_{I}^{*}\left(V_{I}\right)$ in $\operatorname{Comp}_{I}^{*}(M)$.

Definition 4.12 (Compressed multiplicities). Let $M$ be a representation of $\vec{G}_{m, n}$ and $I \in \mathbb{I}_{m, n}$. Define the source-sink (ss)-compressed multiplicity as

$$
\underline{d}_{M}^{\mathrm{ss}}(I):=d_{\operatorname{Comp}_{I}^{\mathrm{ss}}(M)}\left(\operatorname{Comp}_{I}^{\mathrm{ss}}\left(V_{I}\right)\right) .
$$

While not the main focus of this paper, for completeness we also define the compressed multiplicities

$$
\underline{d}_{M}^{\mathrm{cc}}(I):=d_{\operatorname{Comp}_{I}^{\mathrm{cc}}(M)}\left(\operatorname{Comp}_{I}^{\mathrm{cc}}\left(V_{I}\right)\right),
$$

and

$$
\underline{d}_{M}^{\mathrm{tot}}(I):=d_{\operatorname{Comp}_{I}^{\mathrm{tot}}(M)}\left(\operatorname{Comp}_{I}^{\mathrm{tot}}\left(V_{I}\right)\right)
$$

In the above, $d_{?}(-)$ is the usual multiplicity function obtained from Theorem 2.1.
One motivation for the above definitions is that we want to compute the multiplicity of an interval module $V_{I}$ as a direct summand of $M$. However, as this may not be straightforward, we instead compute the multiplicity with respect to compressed versions of $M$ and $I$. The rest of this section is devoted to exploring the consequences of this approach.

Remark 4.13. Let $\operatorname{rk}(M): \operatorname{Con}(P) \rightarrow \mathcal{J}(\mathcal{C})$ be the generalized rank invariant as defined in [14], applied to the setting we consider. That is, $P$ is the poset corresponding to the $m \times n$ commutative grid, and the target category is $\mathcal{J}(\mathcal{C})=\mathcal{J}\left(\right.$ vect $\left._{K}\right)$, the category of isomorphism classes of $K$-vector spaces. By definition
$\operatorname{Con}(P)$ is the set of path-connected subposets of $P$, which contains the set of intervals. See [14] for more detailed definitions. We note that for $I \in \mathbb{I}_{m, n}$, the equality

$$
\underline{d}_{M}^{\mathrm{tot}}(I)=\operatorname{dim} \operatorname{rk}(M)(I)
$$

holds. This follows immediately from Lemma 3.1 of [18] applied to $\operatorname{Comp}_{I}^{\text {tot }}(M)$. That is, for intervals $I$, the tot-compressed multiplicity coincides with the generalized rank invariant of [14].

As the next example shows, the values of $\underline{d}_{M}^{\text {ss }}(I)$ and $\underline{d}_{M}^{\text {tot }}(I)=\operatorname{dim} \operatorname{rk}(M)(I)$ can be different in general.
Example 4.14. Let $M$ be the representation of $\vec{G}_{2,3}$ given by


For the interval

it can be computed that $\underline{d}_{M}^{\text {ss }}(I)=1$ while $\underline{d}_{M}^{\text {tot }}(I)=0$.
Remark 4.15. However, if we allow to change the form of the "input" to the function $\operatorname{dim} \operatorname{rk}(M)(-)$ and broaden its domain of definition, the equality $\underline{d}_{M}^{\text {ss }}(I)=\operatorname{dimrk}(M)(\operatorname{Source}(I) \cup \operatorname{Sink}(I))$ holds by the same reasoning as the previous remark. Note that in general, Source $(I) \cup \operatorname{Sink}(I)$ is not necessarily a pathconnected subposet ([14, Definition 2.16]), and thus the original definition of the generalized rank invariant cannot be used. That is, the values of the source-sink compressed multiplicity can be expressed as some value of the generalized rank invariant suitably generalized.

### 4.3. Rank invariant and dimension vector as compression

In this subsection, we show that the compressed multiplicity generalizes the rank invariant [4], a wellknown invariant for 2 D persistence modules.

Recall that the rank invariant is the function assigning to each pair $s, t \in \vec{G}_{m, n}$ with a path from $s$ to $t$, the value

$$
\operatorname{rank}(M(s \rightarrow t))
$$

where $M(s \rightarrow t): M(s) \rightarrow M(t)$ is the linear map associated by $M$ to a path from $s$ to $t$. Note that this is well-defined due to the commutativity relations imposed on $M$.

An interval $R=\bigsqcup_{i=x}^{y}\left[b_{i}, d_{i}\right]_{i} \in \mathbb{I}_{m, n}$ is said to be a rectangle if $b_{i}=b_{i+1}$ and $d_{i}=d_{i+1}$ for any $i=x, \cdots, y-1$. The set of rectangles in $\mathbb{I}_{m, n}$ is denoted by $R_{m, n}$. It is immediate that any rectangle $R$ has a unique source $s$ and a unique $\operatorname{sink} t$. Below is an example of a rectangle together with its source and sink.


We comment that if $\vec{G}_{m, n}$ is viewed as a subposet of $\mathbb{Z} \times \mathbb{Z}$ with coordinate-wise $\leq$, the rectangle $R$ is in fact the segment $R=[s, t]$ in the poset $\mathbb{Z} \times \mathbb{Z}$. In this work, we do not directly use this point of view since we defined $\vec{G}_{m, n}$ as a bound quiver and not as a poset.

Conversely, given any pair $s, t \in \vec{G}_{m, n}$ with a path from $s$ to $t$ (as in the definition of the rank invariant), there is a unique rectangle $R$ with source $s$ and $\operatorname{sink} t$. Thus, the rank invariant can be equivalently defined as the function assigning to each rectangle $R$ in $\mathbb{I}_{m, n}$ the value $\operatorname{rank}(M(s \rightarrow t))$, where $s$ is the unique source of $R$ and $t$ the unique sink.

Let $R$ be a rectangle with source $s$ and $\operatorname{sink} t$. Let us compute the values of the compressed multiplicities at $R$.

- The ss-compressed category of $R$ is: $R^{\mathrm{ss}}: s \longrightarrow t$, so that $\operatorname{Comp}_{R}^{\mathrm{ss}}(M)$ is $M(s) \xrightarrow{M(s \rightarrow t)} M(t)$. Note that a linear map $f: V \rightarrow W$ between finite-dimensional vector spaces is equivalent to the direct sum $(K \rightarrow 0)^{\operatorname{dim} \operatorname{ker} f} \oplus(K \xrightarrow{1} K)^{\operatorname{rank} f} \oplus(0 \rightarrow K)^{\operatorname{dim} \operatorname{coker} f}$. Then we compute

$$
\begin{aligned}
\underline{d}_{M}^{\mathrm{ss}}(R) & =d_{\left(\operatorname{Comp}_{R}^{\mathrm{ss}(M))}\right.}\left(\operatorname{Comp}_{R}^{\mathrm{ss}}\left(V_{R}\right)\right) \\
& =d(M(s) \xrightarrow{M(s \rightarrow t)} M(t))(K \xrightarrow{1} K) \\
& =\operatorname{rank}(M(s \rightarrow t))
\end{aligned}
$$

- Since $R$ has source $s$ and $\operatorname{sink} t$ together with its two other corners (say $u$ and $w$ ) as its cc-essential vertices, the cc-compressed category of $R$ is:
so that $\operatorname{Comp}_{R}^{\mathrm{cc}}(M)$ is

$$
\begin{gathered}
M(u) \longrightarrow M(t) \\
\uparrow \\
M(s) \longrightarrow M(w)
\end{gathered}
$$

Furthermore, $\operatorname{Comp}_{R}^{\mathrm{cc}}\left(V_{R}\right)$ is the injective indecomposable representation $I(t)$ associated to the vertex $t$ :

$$
I(t)=\begin{array}{r}
\underset{1}{K} \xrightarrow{1} K \\
K \xrightarrow{\mid} K \\
K \xrightarrow{1},
\end{array}
$$

and so by [19, Theorem 3 (see also Example 3)]

$$
\begin{aligned}
& =\operatorname{dim}_{\operatorname{Hom}_{R^{\mathrm{cc}}}(I(t), M)-\operatorname{dim} \operatorname{Hom}_{R^{\mathrm{cc}}}(I(t) / \operatorname{soc} I(t), M)} \\
& =\operatorname{dim} M(s)-(\operatorname{dim} M(s)-\operatorname{rank}(M(s \rightarrow t))) \\
& =\operatorname{rank}(M(s \rightarrow t)) \text {. }
\end{aligned}
$$

In the above, soc $I(t)$ is the socle of $I(t)$, which is the sum of all simple submodules of $I(t)$ by definition.

- Finally, $\operatorname{Comp}_{R}^{\mathrm{tot}}(M)$ is the representation of $R^{\mathrm{tot}}$ obtained by restricting $M$ to the rectangle $R$. Furthermore, $\operatorname{Comp}_{R}^{\text {tot }}\left(V_{R}\right)$ is the injective indecomposable representation $I(t)$ of $R^{\text {tot }}$, and

$$
\underline{d}_{M}^{\mathrm{tot}}(R)=\operatorname{rank}(M(s \rightarrow t))
$$

follows from [19, Theorem 3], as above.
The above considerations prove the following.

Proposition 4.16. Let $M$ be a representation of $\vec{G}_{m, n}$ and $R$ a rectangle. For $*=s s, c c$, tot, we have

$$
\underline{d}_{M}^{*}(R)=\operatorname{rank} M(s \rightarrow t)
$$

where $s$ is the unique source vertex of $R$ and $t$ is the unique sink vertex of $R$.
In this sense, the compressed multiplicities $\underline{d}_{M}^{*}(-)$ are generalizations of the rank invariant. With our invariant we hope to capture finer information that cannot be detected by just the rank invariant.

Next, we give an example of representations with the same rank invariants but different compressed multiplicities for intervals that are not rectangles.
 the following representations of $\vec{G}_{2,2}$ :

Clearly, rank invariants of $M$ and $N$ coincide. However, we have $\underline{d}_{M}^{\mathrm{ss}}(I)=1 \neq 0=\underline{d}_{N}^{\mathrm{ss}}(I)$.
We end this subsection with the following observation.
Proposition 4.18. Let $M$ be a representation of $\vec{G}_{m, n}$ and $i$ a vertex of $\vec{G}_{m, n}$. For $*=s s, c c$, tot, we have

$$
\underline{d}_{M}^{*}(\{i\})=\operatorname{dim} M(i),
$$

where $\{i\}$ is the interval subquiver consisting of only the vertex $i$.
Proof. A direct computation shows that

$$
\underline{d}_{M}^{*}(\{i\})=d_{\operatorname{Comp}_{\{i\}}^{*}(M)}\left(\operatorname{Comp}_{\{i\}}^{*}\left(V_{\{i\}}\right)\right)=d_{M(i)}(K)=\operatorname{dim} M(i)
$$

Alternatively, this follows immediately from Proposition 4.16 by considering the rectangle with $s=t=i$.
Namely, the compressed multiplicities $\underline{d}_{M}^{*}(-)$ restricted to vertices coincide with the dimension vector of $M$.

### 4.4. Compression and Inversion

Next, we derive some basic properties of $\underline{d}_{M}^{*}(-)$, and end this section with Theorem 4.23 , which states that for interval-decomposable representations $M$, we can recover the true multiplicity function $d_{M}$ using $\underline{d}_{M}^{*}(-)$.

First, we start with some Lemmas that lead to a Key Lemma 4.21.
Lemma 4.19. If a representation $M$ of $\vec{G}_{m, n}$ decomposes as $M=M_{1} \oplus M_{2}$, then

$$
\underline{d}_{M}^{*}(I)=\underline{d}_{M_{1}}^{*}(I)+\underline{d}_{M_{2}}^{*}(I)
$$

for $*=s s, c c$, tot.
Proof. Since the compression functor $\operatorname{Comp}_{I}^{*}(-)$ is additive, we have $\operatorname{Comp}_{I}^{*}(M)=\operatorname{Comp}_{I}^{*}\left(M_{1}\right) \oplus \operatorname{Comp}_{I}^{*}\left(M_{2}\right)$. Then the statement follows by the Krull-Schmidt theorem.

Lemma 4.20. Let $I, J$ be intervals of $\vec{G}_{m, n}$. Then

$$
\underline{d}_{V_{J}}^{*}(I)= \begin{cases}1 & \text { if } J \in U(I) \quad(\text { i.e. } I \leq J) \\ 0 & \text { otherwise }\end{cases}
$$

for $*=s s, c c$, tot.
Proof. If $I \leq J$, then $\operatorname{Comp}_{I}^{*}\left(V_{J}\right)=\operatorname{Comp}_{I}^{*}\left(V_{I}\right)$, thus $\underline{d}_{V_{J}}^{*}(I)=1$.
On the other hand, if $I \not \leq J$, then there exists some $i \in I_{0}^{*} \backslash J_{0}$ by Lemma 4.3 or Lemma 4.6 for $*=\mathrm{ss}, \mathrm{cc}$, respectively, and by the fact that $I_{0}^{\text {tot }}=I_{0}$, for $*=$ tot. Thus, $i \in \operatorname{supp}\left(\operatorname{Comp}_{I}^{*}\left(V_{I}\right)\right)$ but $i \notin$ $\operatorname{supp}\left(\operatorname{Comp}_{I}^{*}\left(V_{J}\right)\right)$. This means that $\operatorname{Comp}_{I}^{*}\left(V_{J}\right)$ does not have a direct summand isomorphic to $\operatorname{Comp}_{I}^{*}\left(V_{I}\right)$, showing that $\underline{d}_{V_{J}}^{*}(I)=0$.

Lemma 4.21 (Key Lemma). Let $M$ be an interval-decomposable representation of $\vec{G}_{m, n}$ and $I$ an interval in $\mathbb{I}_{m, n}$. Then

$$
\underline{d}_{M}^{*}(I)=\sum_{J \in U(I)} d_{M}\left(V_{J}\right)
$$

for $*=s s, c c$, tot.
Proof. Let $M \cong \bigoplus_{J \in \mathbb{I}_{m, n}} V_{J}^{d_{M}\left(V_{J}\right)}$ be an interval decomposition of a representation $M$ of $\vec{G}_{m, n}$. Then

$$
\underline{d}_{M}^{*}(I)=\sum_{J \in \mathbb{I}_{m, n}} d_{M}\left(V_{J}\right) \cdot \underline{d}_{V_{J}}^{*}(I)=\sum_{J \in U(I)} d_{M}\left(V_{J}\right)
$$

by Lemmas 4.19 and 4.20 .
As a consequence, in the case that $M$ is interval-decomposable, $\underline{d}_{M}^{*}(I)$ does not depend on $*$.
Readers familiar with the Möbius theory for (locally-finite) posets [20] may recognize that Lemma 4.21 simply states that for interval-decomposable representations, the function $\underline{d}_{M}^{*}(-)$ is equal to $d_{M}(-)$ multiplied by the zeta function. Theorem 4.23 below can then be seen as an application of Möbius inversion. Here, we give a direct proof of Theorem 4.23 and delay these Möbius-theoretic considerations to a later section.

First, we note the following proposition which follows immediately from Lemma 4.21.
Proposition 4.22. Let $M$ be an interval-decomposable representation of $\vec{G}_{m, n}$ and $I$ an interval in $\mathbb{I}_{m, n}$. Then

$$
d_{M}\left(V_{I}\right)=\underline{d}_{M}^{*}(I)-\sum_{J \in U(I) \backslash\{I\}} d_{M}\left(V_{J}\right) .
$$

for $*=s s, c c$, tot.
Theorem 4.23 (For interval-decomposables, compressed multiplicity recovers the multiplicity). Let $M$ be an interval decomposable representation of $\vec{G}_{m, n}$ and $I$ an interval in $\mathbb{I}_{m, n}$. Then:

$$
d_{M}\left(V_{I}\right)=\underline{d}_{M}^{*}(I)+\sum_{\emptyset \neq S \subseteq \operatorname{Cov}(I)}(-1)^{\# S} \underline{d}_{M}^{*}(\bigvee S)
$$

for $*=s s, c c$, tot.
Proof. We define the function $f: 2^{U(I)} \rightarrow \mathbb{Z}$ by $f(S):=\sum_{J \in S} d_{M}\left(V_{J}\right)$ for $S \in 2^{U(I)}$, where $2^{U(I)}$ is the power set of $U(I)$. Rewriting Proposition 4.22, we have

$$
d_{M}\left(V_{I}\right)=\underline{d}_{M}^{*}(I)-f\left(\bigcup_{J \in \operatorname{Cov}(I)} U(J)\right)
$$

since $U(I) \backslash\{I\}=\bigcup_{J \in \operatorname{Cov}(I)} U(J)$. Here, the inclusion-exclusion principle ${ }^{3}$ shows that

$$
f\left(\bigcup_{J \in \operatorname{Cov}(I)} U(J)\right)=\sum_{\emptyset \neq S \subseteq \operatorname{Cov}(I)}(-1)^{(\# S-1)} f\left(\bigcap_{J \in S} U(J)\right)
$$

By Proposition 3.6, the join $\bigvee S$ in $U(I)$ exists, and it can be checked that

$$
\bigcap_{J \in S} U(J)=U(\bigvee S)
$$

by definition. Therefore

$$
f\left(\bigcap_{J \in S} U(J)\right)=f(U(\bigvee S))=\underline{d}_{M}^{*}(\bigvee S)
$$

by Lemma 4.21, which completes our proof.
Theorem 4.23 says that to calculate $d_{M}\left(V_{I}\right)$, it is enough to calculate $\underline{d}_{M}^{\text {ss }}(J)$ (which is equal to $\underline{d}_{M}^{\text {cc }}(J)$ and also to $\underline{d}_{M}^{\text {tot }}(J)$ since $M$ is interval-decomposable) for certain intervals $\bar{J}$. We warn that the assumption that $M$ is interval-decomposable is necessary for Key Lemma 4.21, and so is also necessary here. It is easy to construct examples where the equality in Theorem 4.23 fails for non-interval-decomposable representations.
Example 4.24. Let us follow the proof of Theorem 4.23 by computing a particular example. Let $M$ be an interval-decomposable representation of $\vec{G}_{2,4}$ and let $I=\left(\begin{array}{cccc}0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0\end{array}\right) \in \mathbb{I}_{2,4}$, an interval. In this case,

$$
\operatorname{Cov}(I)=\left\{I_{1}:=\left(\begin{array}{llll}
1 & 1 & 1 & 0 \\
0 & 1 & 1 & 0
\end{array}\right), I_{2}:=\left(\begin{array}{llll}
0 & 1 & 1 & 0 \\
0 & 1 & 1 & 1
\end{array}\right)\right\}
$$

and $I_{1} \vee I_{2}=\left(\begin{array}{lll}1 & 1 & 1 \\ 0 & 1 & 1\end{array}\right)$. By Lemma 4.21, we have

$$
\begin{aligned}
& \underline{d}_{M}^{*}(I)=\sum_{J \in U(I)} d_{M}(J) \\
& =d_{M}\left(\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 1 & 1
\end{array}\right)\right)+d_{M}\left(\left(\begin{array}{llll}
1 & 1 & 0 \\
0 & 1 & 1 & 0
\end{array}\right)\right)+d_{M}\left(\left(\begin{array}{llll}
0 & 1 & 1 & 0 \\
0 & 1 & 1 & 1
\end{array}\right)\right) \\
& +d_{M}\left(\left(\begin{array}{lll}
1 & 1 & 1 \\
1 & 1 & 1
\end{array}\right)\right)+d_{M}\left(\left(\begin{array}{lll}
1 & 1 & 1 \\
0 & 1 & 0 \\
0 & 1 & 1
\end{array}\right)\right)+d_{M}\left(\left(\begin{array}{llll}
0 & 1 & 1 & 1 \\
0 & 1 & 1 & 1
\end{array}\right)\right) \\
& +d_{M}\left(\left(\left(\begin{array}{lll}
1 & 1 & 1 \\
1 & 1 & 0
\end{array}\right)\right)+d_{M}\left(\left(\left(\begin{array}{lll}
0 & 1 & 1 \\
0 & 1 & 1
\end{array}\right)\right)+d_{M}\left(\left(\begin{array}{lll}
1 & 1 & 1
\end{array}\right)\right.\right.\right. \\
& \left.=d_{M}\left(V_{I}\right)+\sum_{J \in\left(U\left(I_{1}\right) \cup U\left(I_{2}\right)\right)} d_{M}(J)\right) \\
& =d_{M}\left(V_{I}\right)+\sum_{J \in U\left(I_{1}\right)} d_{M}(J)+\sum_{J \in U\left(I_{2}\right)} d_{M}(J)-\sum_{J \in\left(U\left(I_{1}\right) \cap U\left(I_{2}\right)\right)} d_{M}(J) \\
& =d_{M}\left(V_{I}\right)+\sum_{J \in U\left(I_{1}\right)} d_{M}(J)+\sum_{J \in U\left(I_{2}\right)} d_{M}(J)-\sum_{J \in U\left(I_{1} \vee I_{2}\right)} d_{M}(J) .
\end{aligned}
$$

We thus have

$$
d_{M}\left(V_{I}\right)=\underline{d}_{M}^{*}(I)-\underline{d}_{M}^{*}\left(I_{1}\right)-\underline{d}_{M}^{*}\left(I_{2}\right)+\underline{d}_{M}^{*}\left(I_{1} \vee I_{2}\right)
$$

which is also given by Theorem 4.23.
As another example, let us consider the equioriented $A_{n}$-type quiver, which can be viewed as $\vec{G}_{1, n}$. In this setting, Theorem 4.23 reduces to the following well-known formula. See for example, [2] and Definition 3.2 of [11].
Corollary 4.25. Let $M \in \operatorname{rep} \vec{G}_{1, n}$. For $\mathbb{I}[i, j]$ an interval representation of $\vec{G}_{1, n}$,

$$
\begin{aligned}
d_{M}(\mathbb{I}[i, j])= & {[\operatorname{rank} M((i-1) \rightarrow(j+1))-\operatorname{rank} M((i-1) \rightarrow j)]-} \\
& {[\operatorname{rank} M(i \rightarrow(j+1))-\operatorname{rank} M(i \rightarrow j)], }
\end{aligned}
$$

where if $i-1$ and/or $j+1$ is not in $\vec{G}_{1, n}$, the corresponding term above is 0 .

[^2]Proof. In $\vec{G}_{1, n}$, it follows immediately from the definition that

$$
\underline{d}_{M}^{*}(\mathbb{I}[i, j])=\operatorname{rank} M(i \rightarrow j)
$$

for $*=\mathrm{ss}, \mathrm{cc}$, tot. Furthermore, $\operatorname{Cov}(\mathbb{I}[i, j])$ contains $\mathbb{I}[i-1, j]$ if $i-1 \in \vec{G}_{1, n}$ and contains $\mathbb{I}[i, j+1]$ if $j+1 \in \vec{G}_{1, n}$, and no other elements.

It is well-known that all representations of $\vec{G}_{1, n}$ are interval-decomposable, and thus Theorem 4.23 is applicable. Thus,

$$
\begin{aligned}
d_{M}(\mathbb{I}[i, j])= & \underline{d}_{M}^{*}(\mathbb{I}[i, j]) \\
& -\underline{d}_{M}^{*}(\mathbb{I}[i-1, j])-\underline{d}_{M}^{*}(\mathbb{I}[i, j+1]) \\
& +\underline{d}_{M}^{*}(\mathbb{I}[i-1, j+1])
\end{aligned}
$$

where if $i-1$ and/or $j+1$ is not in $\vec{G}_{1, n}$, the corresponding term above is 0 . Expanding and rearranging terms gives us the required expression.

We note that the same formula has been obtained by using Auslander-Reiten theory in the paper [19] (Equation (9) of [19]). Our Theorem 4.23 here uses only the local lattice structure of $\mathbb{I}_{m, n}$, and it may be interesting to explore Theorem 4.23 using Auslander-Reiten theory, and more generally, a representationtheoretic perspective.

### 4.5. Restriction to equioriented $2 \times n$ commutative grid

In this subsection, we study the special case of $\vec{G}_{2, n}$, which is the equioriented commutative ladder. In this setting, the compressed categories take on very nice forms.

Proposition 4.26. Let $I \in \mathbb{I}_{2, n}$. The quiver of the ss-compressed category $I^{\text {ss }}$ has one of the following forms:
(1) •,
$(2) \bullet \longrightarrow$,
$(3) \bullet \longrightarrow \bullet \bullet$,
(4) $\bullet \longleftrightarrow \bullet \longrightarrow \bullet$,
$(5) \bullet \longrightarrow \bullet \longleftarrow \bullet \longrightarrow$.
Proof. A direct computation shows this.
Similarly, we have the following.
Proposition 4.27. Let $I \in \mathbb{I}_{2, n}$. The bound quiver of the cc-compressed category $I^{c c}$ has one of the following forms:
(1) •,
(2)
$\bullet \longrightarrow$,
(7)

(3)

(4)

(8)

(5)

(6)

(9)


Proof. It is immediate that there are at most 6 cc-essential vertices, arranged in the shape of (9), for an interval in $\mathbb{I}_{2, n}$. The rest of the forms cover the cases where some of those vertices are not cc-essential in $I$.

For $I \in \mathbb{I}_{2, n}$ with $n \geq 5, I^{\text {tot }}$ is of infinite representation type (see [5, Theorem 1.3] or [9] for example). Therefore, it may be difficult to calculate the values $\underline{d}_{M}^{\text {tot }}(I)$.

On the other hand, Proposition 4.26 and Proposition 4.27 shows that $I^{\mathrm{ss}}$ and $I^{\mathrm{cc}}$ are of finite type for any $I \in \mathbb{I}_{2, n}$. In addition, the Auslander-Reiten quivers for the bound quivers in the lists of Proposition 4.26 and Proposition 4.27 can be calculated explicitly. Thus, it is not difficult to calculate the values $\underline{d}_{M}^{*}(I)$ for $*=\mathrm{ss}$, cc, in the setting of the equioriented $2 \times n$ commutative grid.

We discuss more about computations in Section 6.

## 5. Approximation

In this section, let we discuss how to use the above ideas as an approximation of general 2 D persistence modules in $\operatorname{rep} \vec{G}_{m, n}$ by an interval-decomposable one. First, let us rephrase Theorem 4.23 using the language of Möbius inversion, as discussed in Subsection 2.3, with underlying field $F=\mathbb{R}$.

We can view $d_{M}$ and $\underline{d}_{M}^{*}$ as functions $\mathbb{I}_{m, n} \rightarrow \mathbb{R}$ (taking only nonnegative integer values). For $d_{M}$, this is an abuse of notation, since $d_{M}$ is a function from (isomorphism classes of) all indecomposables, but here we are using the symbol to denote it restricted to the interval representations of $\vec{G}_{m, n}$, identified with the set of intervals $\mathbb{I}_{m, n}$.

In the notation of Subsection 2.3, we have $d_{M}, \underline{d}_{M}^{*} \in \mathbb{R}^{\mathbb{I}_{m, n}}$. Then, the Key Lemma 4.21 states that for $M$ interval-decomposable,

$$
\begin{equation*}
\underline{d}_{M}^{*}=\zeta d_{M} \tag{5.1}
\end{equation*}
$$

where the multiplication of $\zeta$ in Eq. (5.1) is precisely the left action of $I\left(\mathbb{I}_{m, n}\right)$ on $\mathbb{R}^{\mathbb{I}_{m, n}}$. By Möbius inversion (multiplication of $\mu=\zeta^{-1}$ ), we obtain

$$
\begin{equation*}
d_{M}=\mu \underline{d}_{M}^{*} . \tag{5.2}
\end{equation*}
$$

This expresses $d_{M}$ in terms of $\underline{d}_{M}^{*}$, a conclusion similar to the one of Theorem 4.23. Next, we show that the coefficients appearing in Theorem 4.23 gives the values of the Möbius function $\mu([I, J])$ of $\mathbb{I}_{m, n}$.

Definition 5.1. Define the function $\mu^{\prime}: \operatorname{Seg}\left(\mathbb{I}_{m, n}\right) \rightarrow \mathbb{R}$, an element of the incidence algebra $I\left(\mathbb{I}_{m, n}\right)$ by the following.

$$
\mu^{\prime}([I, J])=\left\{\begin{array}{lll}
1 & \text { if } I=J  \tag{5.3}\\
\sum_{\substack{J=\bigvee S \\
\emptyset \neq S \subseteq \operatorname{Cov}(I)}}(-1)^{\# S} & \text { otherwise } .
\end{array}\right.
$$

Note that in the case $I \neq J$ and where there is no $\emptyset \neq S \subseteq \operatorname{Cov}(I)$ such that $J=\bigvee S$, the sum above is empty, and thus $\mu^{\prime}([I, J])=0$. The values of $\mu^{\prime}$ are exactly the coefficients appearing in the formula of Theorem 4.23, from which we immediately get the following Corollary.
Corollary 5.2 (Restatement of Theorem 4.23). Let $M$ be an interval-decomposable representation of $\vec{G}_{m, n}$ and $I$ an interval in $\mathbb{I}_{m, n}$. Then:

$$
d_{M}=\mu^{\prime} \underline{d}_{M}^{*}
$$

for $*=s s, c c$, tot.
Theorem 5.3. Let $\mu^{\prime}$ be as defined in Definition 5.1, and $\mu$ be the Möbius function of the poset $\mathbb{I}_{m, n}$. Then,

$$
\mu=\mu^{\prime}
$$

In particular, Equation (5.3) gives the values of $\mu$.

Proof. Let $I \leq L$ be intervals in $\mathbb{I}_{m, n}$. Below, we compare the values $\mu([I, L])$ and $\mu^{\prime}([I, L])$ by induction on L.

First, let us consider $L$ a cover of $I$ and fix $M=V_{L}$. By Corollary 5.2 and Equation 5.2, we have

$$
\mu^{\prime} \underline{d}_{M}^{*}=\mu \underline{d}_{M}^{*} .
$$

We obtain the following sequence of equations by working on both sides the equation.

$$
\begin{aligned}
\left(\mu^{\prime} \underline{d}_{M}^{*}\right)(I) & =\left(\mu \underline{d}_{M}^{*}\right)(I) \\
\sum_{I \leq J} \mu^{\prime}([I, J]) \underline{d}_{M}^{*}(J) & =\sum_{I \leq J} \mu([I, J]) \underline{d}_{M}^{*}(J) \\
\sum_{I \leq J \leq L} \mu^{\prime}([I, J]) \underline{d}_{M}^{*}(J) & =\sum_{I \leq J \leq L} \mu([I, J]) \underline{d}_{M}^{*}(J) \\
\mu^{\prime}([I, I])+\mu^{\prime}([I, L]) & =\mu([I, I])+\mu([I, L]) \\
1+\mu^{\prime}([I, L]) & =1+\mu([I, L]),
\end{aligned}
$$

where going from the second line to the third line follows by Lemma 4.20. We conclude $\mu^{\prime}([I, L])=\mu([I, L])$ for any $L \in \operatorname{Cov}(I)$.

Next, we assume that for any interval $L^{\prime}$ with $L^{\prime}<L, \mu^{\prime}\left(\left[I, L^{\prime}\right]\right)=\mu\left(\left[I, L^{\prime}\right]\right)$. Then we have the following sequence of equations by taking $M=V_{L}$ and again using Lemma 4.20:

$$
\begin{aligned}
\left(\mu^{\prime} \dot{d}_{M}^{*}\right)(I) & =\left(\mu d_{M}^{*}\right)(I) \\
\sum_{I \leq J \leq L} \mu^{\prime}([I, J]) & =\sum_{I \leq J \leq L} \mu([I, J]) \\
\sum_{I \leq J<L} \mu^{\prime}([I, J])+\mu^{\prime}([I, L]) & =\sum_{I \leq J<L} \mu([I, J])+\mu([I, L]) .
\end{aligned}
$$

Since we have $\sum_{I \leq J<L} \mu^{\prime}([I, J])=\sum_{I \leq J<L} \mu([I, J])$ by the inductive assumption, we obtain $\mu^{\prime}([I, L])=\mu([I, L])$. By the induction, we get the conclusion.

As we have seen, $d_{M}=\mu \underline{d}_{M}^{*}$ for $M$ interval-decomposable. Even in the case where $M$ is not intervaldecomposable, we nevertheless can do the transformation. Thus we define $\delta_{M}^{*}:=\mu \underline{d}_{M}^{*}$ in general.
Definition 5.4. Put $*=\mathrm{ss}$, cc, tot. Define $\delta_{M}^{*}:=\mu \underline{d}_{M}^{*}$. In particular, for each $I \in \mathbb{I}_{m, n}$ an interval subquiver of $\vec{G}_{m, n}$,

$$
\delta_{M}^{*}(I):=\underline{d}_{M}^{*}(I)+\sum_{\emptyset \neq S \subseteq \operatorname{Cov}(I)}(-1)^{\# S} \underline{d}_{M}^{*}(\bigvee S)
$$

First, we note the following obvious property of $\delta_{M}^{*}(-)$.
Lemma 5.5. If $M \cong M_{1} \oplus M_{2}$, then we have

$$
\delta_{M}^{*}(-)=\delta_{M_{1}}^{*}(-)+\delta_{M_{2}}^{*}(-)
$$

Proof. Since $\underline{d}_{M}^{*}(-)=\underline{d}_{M_{1}}^{*}(-)+\underline{d}_{M_{2}}^{*}(-)$ by Lemma 4.19, we have the desired equation by definition.
Since in general

$$
M \cong \bigoplus_{X \in \mathcal{L}} X^{d_{M}(X)}
$$

by Theorem 2.1 (where $d_{M}$ is the actual multiplicity function, not restricted to intervals), one way of constructing an approximating interval-decomposable object is to naively define

$$
\begin{equation*}
\delta^{*}(M)=\bigoplus_{I \in \mathbb{I}_{m, n}}\left(V_{I}\right)^{\delta_{M}^{*}(I)} \tag{5.4}
\end{equation*}
$$

by taking the function $\delta_{M}^{*}$ on $\mathbb{I}_{m, n}$ as a substitute for the function $d_{M}$ on $\mathcal{L}$. Defined this way, $M \cong \delta^{*}(M)$ for interval-decomposable $M$. However, the value $\delta_{M}^{*}(I)$ can be negative in general, and thus the direct sum in Eq. (5.4) does not make sense.

For example, we have the following.

Example 5.6. Let $M$ be the representation of $\vec{G}_{2,3}$ given by


The value of $\delta_{M}^{\mathrm{ss}}(I)$ is 0 except in the cases of $I$ being one of the intervals $I_{1}, I_{2}, I_{3}, I_{4}$ given below.
(1) For $I_{1}: \longrightarrow \longrightarrow_{\bullet}^{\bullet}, \delta_{M}^{\mathrm{ss}}\left(I_{1}\right)=-1$,
(2) For $I_{2}$ :

(3) For $I_{3}$ :

(4) For $I_{4}$ : ,$\delta_{M}^{\mathrm{ss}}\left(I_{4}\right)=1$.

Proof. We directly use Definition 5.4 to compute $\delta_{M}^{\text {ss }}\left(I_{1}\right)$. We let $\operatorname{Cov}\left(I_{1}\right)=\left\{I_{2}, I_{5}\right\}$, and let $I_{6}=I_{2} \vee I_{5}$, where the intervals are given below. We first compute the value of the compressed multiplicity $\underline{d}_{M}^{\text {ss }}(-)$ of these intervals. We have:


Thus, by definition,

$$
\delta_{M}^{\mathrm{ss}}\left(I_{1}\right)=0-1-0+0=-1 .
$$

The other computations follow similarly.
For $M$ interval-decomposable, it is clear from the above that all values of $\delta_{M}^{*}$ are nonnegative, as it is equal to $d_{M}$ itself. In the next example we see that the converse does not hold, and so we cannot use the nonnegativity of $\delta_{M}^{*}$ to check for interval-decomposability.
Example 5.7 (Continuation of Example 5.6). There exist $N$ such that $\delta_{N}^{*}$ is nonnegative, but $N$ is not interval-decomposable.

In particular, let $M$ and $I_{i}(i=1,2,3,4)$ be as given in Example 5.6. Then $N:=M \oplus I_{1}$ is such an example.

Proof. Since $N=M \oplus I_{1}, \delta_{N}^{\mathrm{ss}}=\delta_{M}^{\mathrm{ss}}+\delta_{I_{1}}^{\mathrm{ss}}$ by Lemma 5.5. Then we have

$$
\delta_{N}^{\mathrm{ss}}\left(I_{1}\right)=-1+1=0
$$

and $\delta_{N}^{\mathrm{ss}}(I)=\delta_{M}^{\mathrm{ss}}(I)+0 \geq 0$ for all intervals $I \neq I_{1}$. Thus, $\delta_{N}^{\mathrm{ss}}$ is nonnegative, but $N$ is not intervaldecomposable since $M$ is an indecomposable summand of $N$ that is not isomorphic to an interval representation.

To deal with the possibility of negative terms in $\delta_{M}^{*}$ in general, we use the formalism of the split Grothendieck group to express the addition of a negative number of copies of an interval in a direct sum. For more details, see for example the notes [21, Chapter 2].

Definition 5.8. The split Grothendieck group $\operatorname{Gr}(\mathcal{C})$ of an additive category $\mathcal{C}$ is the free abelian group generated by isomorphism classes $[C]$ of objects in $\mathcal{C}$ modulo the relations $\left[C_{1} \oplus C_{2}\right]=\left[C_{1}\right]+\left[C_{2}\right]$ for all objects $C_{1}, C_{2}$ of $\mathcal{C}$. For an object $C$ of $\mathcal{C}$, we denote by $\llbracket C \rrbracket$ the element of $\operatorname{Gr}(\mathcal{C})$ represented by $[C]$.

In the following we consider the split Grothendieck group $\operatorname{Gr}\left(\operatorname{rep} \vec{G}_{m, n}\right)$ of rep $\vec{G}_{m, n}$. Then by the KrullSchmidt theorem we easily see that it has a basis $\{\llbracket L \rrbracket \mid L \in \mathcal{L}\}$, where $\mathcal{L}$ is a complete set of representatives of the isomorphism classes of indecomposable representations of $\vec{G}_{m, n}$ (see [21, Theorem 2.3.6]). Thus each $X \in \operatorname{Gr}\left(\operatorname{rep} \vec{G}_{m, n}\right)$ is uniquely expressed in the form

$$
X=\sum_{L \in \mathcal{L}} a_{L} \llbracket L \rrbracket
$$

with $a_{L} \in \mathbb{Z}$ for all $L \in \mathcal{L}$. Here we define the representations

$$
\begin{equation*}
X_{+}:=\bigoplus_{\substack{L \in \mathcal{L} \\ a_{L} \geq 0}} L^{a_{L}} \quad \text { and } \quad X_{-}:=\bigoplus_{\substack{L \in \mathcal{L} \\ a_{L}<0}} L^{\left(-a_{L}\right)} \tag{5.5}
\end{equation*}
$$

which are called the positive part and the negative part of $X$, respectively. Note that they are representations of $\vec{G}_{m, n}$ with the property that $X=\llbracket X_{+} \rrbracket-\llbracket X_{-} \rrbracket$ because $\llbracket X_{+} \rrbracket=\sum_{\substack{L \in \mathcal{L} \\ a_{L} \geq 0}} a_{L} \llbracket L \rrbracket$ and $\llbracket X_{-} \rrbracket=$ $\sum_{\substack{L \in \mathcal{L} \\ a_{L}<0}}\left(-a_{L}\right) \llbracket L \rrbracket$. Therefore, $X$ can be uniquely presented by the pair $\left(X_{+}, X_{-}\right)$of representations of $\vec{G}_{m, n}$.
Definition 5.9 (interval-decomposable approximation). Let $M \in \operatorname{rep} \vec{G}_{m, n}$. Define the interval-decomposable approximation $\delta^{*}(M)$ of $M$ by

$$
\begin{equation*}
\delta^{*}(M):=\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \llbracket V_{I} \rrbracket \in \operatorname{Gr}\left(\operatorname{rep} \vec{G}_{m, n}\right) \tag{5.6}
\end{equation*}
$$

for $*=\mathrm{ss}, \mathrm{cc}$, or, tot.
By the above observation, $\delta^{*}(M)$ can be expressed by the pair $\left(\delta^{*}(M)_{+}, \delta^{*}(M)_{-}\right)$of interval-decomposable representations, where

$$
\delta^{*}(M)_{+}=\bigoplus_{\substack{I \in \mathbb{I}_{m, n} \\ \delta_{M}(I)>0}} V_{I}^{\delta_{M}^{*}(I)} \text { and } \delta^{*}(M)_{-}=\bigoplus_{\substack{I \in \mathbb{I}_{m, n} \\ \delta_{M}^{*}(I)<0}} V_{I}\left(-\delta_{M}^{*}(I)\right)
$$

Theorem 5.10. Let $M \in \operatorname{rep} \vec{G}_{m, n}$ be interval-decomposable. Then, $\delta^{*}(M)=\llbracket M \rrbracket$, or equivalently, $\delta^{*}(M)_{+} \cong M$ and $\delta^{*}(M)_{-}=0$.
Proof. Because $M$ is interval-decomposable, $\delta_{M}^{*}=d_{M}$. The conclusion follows immediately from this.
Note that the converse trivially holds. If $\delta^{*}(M)=\llbracket M \rrbracket$ then $M$ is interval-decomposable.
Let us discuss the relationship between $M$ and $\delta^{*}(M)$. In particular, we focus on dimension vectors and rank invariants.

Example 5.11 (Continuation of Example 5.6). With the same notation as in Example 5.6, we have the equality

$$
\begin{aligned}
\sum_{I \in \mathbb{I}_{2,3}} \delta_{M}^{\mathrm{ss}}(I) \cdot \underline{\operatorname{dim}}\left(V_{I}\right) & =\left(\begin{array}{lll}
1 & 1 & 1 \\
0 & 1 & 1
\end{array}\right)+\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 1 & 1
\end{array}\right)+\left(\begin{array}{lll}
1 & 1 & 0 \\
0 & 0 & 0
\end{array}\right)-\left(\begin{array}{lll}
1 & 1 & 0 \\
0 & 1 & 1
\end{array}\right) \\
& =\left(\begin{array}{lll}
1 & 2 & 1 \\
0 & 1 & 1
\end{array}\right) \\
& =\underline{\operatorname{dim}}(M) .
\end{aligned}
$$

For $\delta_{M}^{\mathrm{cc}}$, we have a similar equality of the dimension vectors for the example above. This is not a coincidence, and in fact the equality always holds (see Corollary 5.14). First we prove the following stronger statement.
Theorem 5.12. Let $M$ be a representation of $\vec{G}_{m, n}=(Q, R)$, and let $i$ and $j$ be vertices of $Q$ such that there exists a path from $i$ to $j$ in $Q$. Then we have

$$
\begin{equation*}
\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \cdot \operatorname{rank} V_{I}(i \rightarrow j)=\operatorname{rank} M(i \rightarrow j) \tag{5.7}
\end{equation*}
$$

for $*=s s, c c$, tot.
To prove the theorem above we need the following lemma, which is the essence of Theorem 5.12.
Lemma 5.13. Let $M \in \operatorname{rep} \vec{G}_{m, n}$ and $I \in \mathbb{I}_{m, n}$. Then

$$
\underline{d}_{M}^{*}(I)=\sum_{I \leq J \in \mathbb{I}_{m, n}} \delta_{M}^{*}(J)
$$

Proof. This follows from Möbius inversion. That is, by definition $\delta_{M}^{*}:=\mu \underline{d}_{M}^{*}$ and thus

$$
\underline{d}_{M}^{*}=\zeta \delta_{M}^{*}
$$

since $\mu^{-1}=\zeta$. The right-hand side expanded out gives the result.
Then we prove Theorem 5.12.
Proof of Theorem 5.12. Since there is a path from $i$ to $j$, the rectangle with source $i$ and sink $j$ exists. We denote this rectangle with source $i$ and $\operatorname{sink} j$ by $R_{i, j}$.

We note that for an interval $I \in \mathbb{I}_{m, n}$, rank $V_{I}(i \rightarrow j)$ is 1 if and only if $I$ contains the rectangle $R_{i, j}$ and is 0 otherwise. This gives the first equality in the following computation. We have

$$
\begin{aligned}
\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \cdot \operatorname{rank} V_{I}(i \rightarrow j) & =\sum_{R_{i, j} \leq I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \\
& =\underline{d}_{M}^{*}\left(R_{i, j}\right) \\
& =\operatorname{rank} M(i \rightarrow j)
\end{aligned}
$$

where the second equality follows from Lemma 5.13, and the last equality follows by applying Proposition 4.16.

As a corollary of Theorem 5.12, we have the following desired equation for dimension vectors.
Corollary 5.14. Let $M$ be a representation of $\vec{G}_{m, n}$. Then we have

$$
\begin{equation*}
\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \cdot \underline{\operatorname{dim}}\left(V_{I}\right)=\underline{\operatorname{dim}}(M) \tag{5.8}
\end{equation*}
$$

Proof. It is enough to show that for any $i \in G_{0}$,

$$
\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \cdot\left(\underline{\operatorname{dim}}\left(V_{I}\right)\right)_{i}=(\underline{\operatorname{dim}}(M))_{i} .
$$

Note that $\left(\underline{\operatorname{dim}}\left(V_{I}\right)\right)_{i}=\operatorname{rank} V_{I}(i \rightarrow i)$ and $(\underline{\operatorname{dim}}(M))_{i}=\operatorname{rank} M(i \rightarrow i)$, where the path $i \rightarrow i$ means the path $e_{i}$ of length 0 at $i$. Thus, by Theorem 5.12, we obtain the above equation.

Let us give another consequence of this result, which warns us against thinking of approximation in terms of functions.

In general, $M \in \operatorname{rep} \vec{G}_{m, n}$ can be written as $M \cong M_{I} \oplus X$, where $M_{I}$ is interval-decomposable, and $0 \neq X$ has no interval representation as a summand. By Lemma 5.5,

$$
\begin{equation*}
\delta_{M}^{*}=\delta_{M_{I}}^{*}+\delta_{X}^{*}=d_{M_{I}}+\delta_{X}^{*}: \mathbb{I}_{m, n} \rightarrow \mathbb{R} \tag{5.9}
\end{equation*}
$$

where we also use the fact that $\delta_{M_{I}}^{*}=d_{M_{I}}$ because $M_{I}$ is interval-decomposable. Restricted to $\mathbb{I}_{m, n}, d_{M}$ has the same values as $d_{M_{I}}$. Precisely speaking, by our abuse of notation $d_{M}: \mathbb{I}_{m, n} \rightarrow \mathbb{R}$ above is the full multiplicity function $d_{M}$ restricted to the set of interval representations, which can be identified with $\mathbb{I}_{m, n}$. Thus, we may be tempted to think of using $\delta_{M}^{*}$ to approximate $d_{M_{I}}=d_{M}$ as functions on $\mathbb{I}_{m, n}$. To measure the error involved, we use the $\ell_{1}$-norm of functions $f: \mathbb{I}_{m, n} \rightarrow \mathbb{R}$ defined by $\|f\|_{1}=\sum_{I \in \mathbb{I}_{m, n}}|f(I)|$. Let us consider the value of

$$
\left\|\delta_{X}^{*}\right\|_{1}=\left\|\delta_{M}^{*}-d_{M}\right\|_{1} .
$$

We remind the reader that we are considering $d_{M}$ as a function on $\mathbb{I}_{m, n}$ by restriction.
Corollary 5.15. Let $\vec{G}_{m, n}$ be an equioriented commutative grid of size at least $2 \times 5$ or $5 \times 2$. For any $\ell \in \mathbb{N}$, there exists an indecomposable non-interval representation $X \in \operatorname{rep} \vec{G}_{m, n}$, such that

$$
\left\|\delta_{X}^{*}\right\|_{1} \geq \ell
$$

Proof. The construction in [10] provides such an indecomposable non-interval $X \in \operatorname{rep} \vec{G}_{2,5}$ (for $\vec{G}_{m, n}$ larger than $2 \times 5$, we simply pad with zero spaces and zero maps):
where each $E$ is an $\ell \times \ell$ identity matrix, and $J$ is the $\ell \times \ell$ Jordan block with eigenvalue $\lambda=1$.
Let $i$ be one of the vertices such that $X(i)$ has dimension at least $\ell$. We compute:

$$
\begin{aligned}
\ell \leq \operatorname{dim} X(i) & =\sum_{I \in \mathbb{I}_{m, n}} \delta_{X}^{*}(I) \cdot\left(\underline{\operatorname{dim}}\left(V_{I}\right)\right)_{i} \\
& =\sum_{I: i \in I} \delta_{X}^{*}(I) \\
& \leq \sum_{I: i \in I}\left|\delta_{X}^{*}(I)\right| \\
& \leq \sum_{I \in \mathbb{I}_{m, n}}\left|\delta_{X}^{*}(I)\right| \\
& =\left\|\delta_{X}^{*}\right\|_{1},
\end{aligned}
$$

where the first line follows from Corollary 5.14.
Remark 5.16. A simpler proof can be provided, if we allow $X$ to not be indecomposable in the preceeding corollary, as follows. Let $N$ be an indecomposable non-interval representation, which is known to exist. For example, the above indecomposable can be reused. Then, defining $X$ as the direct sum of $\ell \operatorname{copies}$ of $N$, we have that $X$ and

$$
\left\|\delta_{X}^{*}\right\|_{1}=\left\|\sum_{i=1}^{\ell} \delta_{N}^{*}\right\|_{1}=\ell\left\|\delta_{N}^{*}\right\|_{1} \geq \ell
$$

since $\left\|\delta_{N}^{*}\right\|_{1} \geq 1$ (otherwise $\delta_{N}^{*}=0$ and thus $N=0$, a contradiction).

In other words, the "error term" $\left\|\delta_{X}^{*}\right\|_{1}$ can be made arbitrarily large by varying $M$. However, in the above analysis, we considered the "error term" $\left\|\delta_{X}^{*}\right\|_{1}=\left\|\delta_{M}^{*}-d_{M}\right\|_{1}$ where $d_{M}$ is considered as a function on $\mathbb{I}_{m, n}$ by restriction. That is, its values on non-intervals are ignored. A more comprehensive analysis could potentially take into account those terms as well.

Finally, let us give an interpretation of Theorem 5.12 and Corollary 5.14. The left-hand side

$$
\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \cdot \operatorname{rank} V_{I}(i \rightarrow j)
$$

of Equation (5.7) in Theorem 5.12 and the left-hand side

$$
\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \cdot \underline{\operatorname{dim}}\left(V_{I}\right)
$$

of Equation (5.8) in Corollary 5.14 can be viewed as the rank invariant and the dimension vector of the interval-decomposable approximation

$$
\delta^{*}(M)=\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \llbracket V_{I} \rrbracket
$$

respectively. That is, the rank invariant (dimension vector, respectively) of $\delta^{*}(M)$ can be defined by adding the rank invariants (dimension vectors, respectively) of its summands. With this, Theorem 5.12 and Corollary 5.14 simply states that the interval-decomposable approximation $\delta^{*}(M)$ preserves the rank invariant and dimension vector of $M$. It is in this sense that we think of approximating $M$ by $\delta^{*}(M)$.

## 6. Algorithms for equioriented commutative ladders

Let $M$ be a persistence module over an equioriented $m \times n$ commutative grid. For completeness, we first present a high-level overview of an algorithm for the computation of our proposed interval-decomposable approximation $\delta^{*}(M)$. Afterwards, we consider the case of persistence modules over equioriented commutative ladders ( $2 \times n$ commutative grids).

The computation of $\delta^{*}(M)=\sum_{I \in \mathbb{I}_{m, n}} \delta_{M}^{*}(I) \llbracket V_{I} \rrbracket$ of $M$ involves two major steps:
(1) (Algorithm 1) computation of the compressed multiplicity function $\underline{d}_{M}^{*}: \mathbb{I}_{m, n} \rightarrow \mathbb{N}$, defined by

$$
\underline{d}_{M}^{*}(I):=d_{\operatorname{Comp}_{I}^{*}(M)}\left(\operatorname{Comp}_{I}^{*}\left(V_{I}\right)\right)
$$

for $I \in \mathbb{I}_{m, n}$, and
(2) (Algorithm 2) computation of the Möbius inversion $\delta_{M}^{*}=\mu \underline{d}_{M}^{*}$ given by

$$
\delta_{M}^{*}(I):=\underline{d}_{M}^{*}(I)+\sum_{\emptyset \neq S \subseteq \operatorname{Cov}(I)}(-1)^{\# S} \underline{d}_{M}^{*}(\bigvee S)
$$

for $I \in \mathbb{I}_{m, n}$.
Algorithm 1 below for the computation of the compressed multiplicity simply expands upon the definition.

```
Algorithm 1 Compressed multiplicity \(\underline{d}_{M}^{*}\) of \(M\)
    function CompressedMultiplicity \((M)\)
        Initialize the function \(\underline{d}_{M}^{*}\) on \(\mathbb{I}_{m, n}\)
        for \(I \in \mathbb{I}_{m, n}\) do
            Compute the compressed representation \(M^{\prime}=\operatorname{Comp}_{I}^{*}(M)\).
            Compute the compressed representation \(I^{\prime}=\operatorname{Comp}_{I}^{*}\left(V_{I}\right)\).
                (which is simply the interval representation with the whole of \(I^{*}\) as support)
            Compute the multiplicity \(d_{M^{\prime}}\left(I^{\prime}\right)\) of \(I^{\prime}\) in \(M^{\prime}\).
            \(\underline{d}_{M}^{*}(I) \leftarrow d_{M^{\prime}}\left(I^{\prime}\right)\)
        end for
        return \(\underline{d}_{M}^{*}\)
    end function
```

Line 4 of Algorithm 1 for the compressed representation $M^{\prime}=\operatorname{Comp}_{I}^{*}(M)$ simply means forgetting about the vector spaces (internal linear maps, resp.) of $M$ corresponding to objects (morphisms, resp.) not in the compressed category $I^{*}$. Note that depending on how $M$ is stored, extra computations are needed (if some of the internal maps of $M$ are not explicitly stored, they may need to be computed explicitly and stored if they rely on internal maps about to be forgotten). We provide an example of this with the $2 \times n$ case later.

In general, the computation of the multiplicity $d_{M^{\prime}}\left(I^{\prime}\right)$ of $I^{\prime}$ in $M^{\prime}$ (Line 6 of Algorithm 1) can be accomplished by computing the dimensions of certain homomorphism spaces to entries in the almost split sequence ${ }^{4}$ starting at $I^{\prime}$ (see [19, Theorem 3], [22, Corollary. 2.3] and also [8, Algorithms 3, 4]).

Algorithm 2 is also a straightforward expansion of the definition.

```
Algorithm 2 Möbius inversion \(\delta_{M}^{*}\) of \(\underline{d}_{M}^{*}\)
    function MöbiusInversion \(\left(\underline{d}_{M}^{*}\right)\)
        Initialize the function \(\delta_{M}^{*}\) on \(\mathbb{I}_{m, n}\)
        for \(I \in \mathbb{I}_{m, n}\) do
            \(a \leftarrow \underline{d}_{M}^{*}(I)\)
            Compute \(\operatorname{Cov}(I)\)
            for \(\emptyset \neq S \subseteq \operatorname{Cov}(I)\) do
                Compute \(\bigvee S\)
                \(a \leftarrow a+(-1)^{\# S} \underline{d}_{M}^{*}(\bigvee S)\)
            end for
            \(\delta_{M}^{*}(I) \leftarrow a\)
        end for
        return \(\delta_{M}^{*}\)
    end function
```

Algorithm 2 requires the computation of joins of cover elements of $I$. We comment on this below. Let $I=\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i}$. By Proposition 3.4, the elements of $\operatorname{Cov}(I)$ are given by a specific form. We recall that Proposition 3.4 only provides a list of candidates, from which picking up all valid intervals forms $\operatorname{Cov}(I)$. We single out the following four potential cover elements specified by Proposition 3.4 that need special consideration:

[^3](1) extension of the last row of $I$ by one adjacent vertex left of the row (top-left)
\[

C_{t l}=\bigsqcup_{i=s}^{t}\left[b_{i}^{\prime}, d_{i}\right]_{i}, where b_{i}^{\prime}= $$
\begin{cases}b_{i}-1 & \text { if } i=t \\ b_{i} & \text { otherwise }\end{cases}
$$
\]

(2) extension of the first row of $I$ by one adjacent vertex right of the row (bottom-right)

$$
C_{b r}=\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}^{\prime}\right]_{i}, \text { where } d_{i}^{\prime}= \begin{cases}d_{i}+1 & \text { if } i=s \\ d_{i} & \text { otherwise }\end{cases}
$$

(3) addition of one vertex above the upper-left vertex of $I$ (top)

$$
C_{t}=\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i} \sqcup\left[b_{t}, b_{t}\right]_{t+1},
$$

(4) addition of one vertex below the lower-right vertex of $I$ (bottom)

$$
C_{b}=\left[d_{s}, d_{s}\right]_{s-1} \sqcup \bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i} .
$$

Remark 6.1. Then, it is clear that if $S \subset \operatorname{Cov}(I)$

- does not contain both $C_{t l}$ and $C_{t}$, and
- does not contain both $C_{b r}$ and $C_{b}$,
then $\bigvee S=\bigcup_{C \in S} C$. That is, simply taking the union is enough since the union is an interval.
Otherwise, we need to add at most two vertices to $\bigcup_{C \in S} C$ in order to obtain $\bigvee S$. If $S \subset \operatorname{Cov}(I)$ contains both $C_{t l}$ and $C_{t}$, then an additional vertex in the top left needs to be added to form an interval. Similarly, if $S \subset \operatorname{Cov}(I)$ contains both $C_{b r}$ and $C_{b}$, then an additional vertex in the bottom right needs to be added to form an interval.
Example 6.2. We provide an example using the interval $I$ in the commutative grid $\vec{G}_{5,6}$ with candidate vertices marked as in Example 3.5.


In dimension vector notation,

$$
I=\left(\begin{array}{llll}
0 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0
\end{array}\right)
$$

and all the cover elements are given by

$$
\left.\begin{array}{rl}
C_{t l} & =\left(\begin{array}{llll}
0 & 0 & 0 & 0 \\
1 & 1 & 1 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0
\end{array}\right), C_{t}=\left(\begin{array}{llll}
0 & 1 & 0 & 0 \\
0 & 1 & 1 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0
\end{array}\right), \\
C_{b r} & =\left(\begin{array}{lll}
0 & 0 & 0
\end{array}\right) \\
0 & 1
\end{array}\right)
$$

Thus, for example,

$$
C_{b r} \vee C_{t} \vee C=\left(\begin{array}{ccccc}
0 & 1 & 0 & 0 \\
0 & 1 & 1 & 0 \\
0 & 1 & 1 & 1 \\
0 & 0 & 0 & 0
\end{array}\right)=C_{b r} \cup C_{t} \cup C
$$

while

$$
C_{t l} \vee C_{t} \vee C_{b r}=\left(\begin{array}{ccc}
1 & 1 & 0
\end{array}\right)
$$

where $v$ is the vertex at the upper-left corner.
Theorem 6.3. Algorithm 2, which computes $\delta_{M}^{*}$ given $\underline{d}_{M}^{*}$, can be performed with time complexity $O\left(\# \mathbb{I}_{m, n} 2^{D} D m\right)$, where $D=\max _{I \in \mathbb{I}_{m, n}} \# \operatorname{Cov}(I)$.

Proof. For each $I \in \mathbb{I}_{m, n}$, there are at most $2^{D}-1$ nonempty subsets $S$ of $\operatorname{Cov}(I)$. By the formula

$$
\delta_{M}^{*}(I):=\underline{d}_{M}^{*}(I)+\sum_{\emptyset \neq S \subseteq \operatorname{Cov}(I)}(-1)^{\# S} \underline{d}_{M}^{*}(\bigvee S)
$$

for each $S$, we need to compute $\bigvee S$, which is the join of at most $D$ intervals.
We first compute $\bigcup_{C \in S} C$ by the following. Assuming that intervals are represented in the form of $I=$ $\bigsqcup_{i=s}^{t}\left[b_{i}, d_{i}\right]_{i}$ (row-wise), with the number of rows equal to $m$, the union of two cover elements can be computed by iterating through the $m$ rows and taking the union of the corresponding intervals $\left[b_{i}, d_{i}\right]_{i} \cup\left[b_{i}^{\prime}, d_{i}^{\prime}\right]_{i}$. We iterate over the elements of $S$ (at most $D$ ) to obtain $\bigcup_{C \in S} C$.

Finally, the above discussion around Remark 6.1 concerning the four cover elements $C_{t l}, C_{t}, C_{b r}, C_{b}$ that need special consideration provides the computation of $\bigvee S$ by modifying the union $\bigcup_{C \in S} C$. We simply need to check for the presence of both $C_{t l}$ and $C_{t}$ in $S$, and both $C_{b r}$ and $C_{b}$ in $S$, and add the additional vertices to $\bigcup_{C \in S} C$ to obtain $\bigvee S$, as noted in Remark 6.1.

By the above, we have as an upper bound $\# \mathbb{I}_{m, n} \cdot\left(2^{D}-1\right) \cdot D \cdot m$ operations, giving the claimed time complexity.

Next, we consider the case of equioriented commutative ladders with $*=$ ss, where it has been noted in Subsection 4.5 that the ss-compressed category is of Dynkin $A_{n}$-type with $n \leq 4$ (Proposition 4.26). So, let $M$ be a persistence module over the $2 \times n$ commutative grid, and let

$$
d=\max _{v \in\left(\vec{G}_{2, n}\right)_{0}} \operatorname{dim} M(v) .
$$

In particular $M \in \operatorname{rep} \vec{G}_{2, n}$ is given as the following collection of vector spaces and linear maps

$$
\begin{aligned}
& M(2,1) \xrightarrow{M((2,1) \rightarrow(2,2))} M(2,2) \xrightarrow{M((2,1) \rightarrow(2,3))} \cdots \xrightarrow{M((2, n-1) \rightarrow(2, n))} M(2, n) \\
& \quad \uparrow M((1,1) \rightarrow(2,1)) \\
& \uparrow M((1,2) \rightarrow(2,2)) \\
& M(1,1) \xrightarrow{M((1,1) \rightarrow(1,2))} M(1,2) \xrightarrow{M((1,2) \rightarrow(1,3))} \cdots \xrightarrow{M((1, n) \rightarrow(2, n)) \uparrow} \\
& M(1, n-1) \rightarrow(1, n)) \\
& M(1, n)
\end{aligned}
$$

such that

$$
\begin{gathered}
M(2, j) \xrightarrow{M((2, j) \rightarrow(2, j+1))} M(2, j+1) \\
M((1, j) \rightarrow(2, j)) \uparrow \\
M(1, j) \xrightarrow{M((1, j) \rightarrow(1, j+1))} M(1, j+1)
\end{gathered}
$$

commutes for all $j \in\{1,2, \ldots, n-1\}$. For $(x, y),(i, j)$ distinct vertices of $\vec{G}_{2, n}$ such that $x \leq i$ and $y \leq j$ (that is, there exists a path from $(x, y)$ to $(i, j)$ in $\left.\vec{G}_{2, n}\right), M((x, y) \rightarrow(i, j))$ is the composition $M(p)=M\left(\alpha_{\ell}\right) \cdots M\left(\alpha_{1}\right)$ where $p=\left(\alpha_{l}, \ldots, \alpha_{\ell}\right)$ is a path from $(x, y)$ to $(i, j)$ in $\vec{G}_{2, n}$. Note that by the commutativity relations, the composition does not depend on the path chosen.

In Algorithm 3, we specialize Algorithm 1 to this setting and add more details. In particular, we precompute all the compositions $M((x, y) \rightarrow(i, j))$ (as each will be used at some point in the algorithm, anyway), and explicitly write down formulae for $d_{M^{\prime}}\left(I^{\prime}\right)$ using ranks of certain matrices.

```
Algorithm 3 ss-compressed multiplicity ( \(2 \times n\) case)
    function SSCompressedMultiplicityTwoByN \((M)\)
        Initialize the function \(\underline{d}_{M}^{\mathrm{ss}}\) on \(\mathbb{I}_{2, n}\)
        Compute \(M((x, y) \rightarrow(i, j))\) for all \((x, y) \neq(i, j)\) with a path from \((x, y)\) to \((i, j)\)
        for \(I \in \mathbb{I}_{2, n}\) do
            Compute \(d_{M^{\prime}}\left(I^{\prime}\right)\) using the formula in Proposition 6.4,
                where \(M^{\prime}=\operatorname{Comp}_{I}^{\mathrm{ss}}(M)\) and \(I^{\prime}=\operatorname{Comp}_{I}^{\mathrm{ss}}\left(V_{I}\right)\).
            \(\underline{d}_{M}^{\mathrm{ss}}(I) \leftarrow d_{M^{\prime}}\left(I^{\prime}\right)\)
        end for
        return \(\underline{d}_{M}^{\text {ss }}\)
    end function
```

Proposition 6.4. Let $M \in \operatorname{rep} K \vec{G}_{2, n}, I \in \mathbb{I}_{2, n}$ and let $M^{\prime}=\operatorname{Comp}_{I}^{s s}(M)$ and $I^{\prime}=\operatorname{Comp}_{I}^{s s}\left(V_{I}\right)$ be their respective compressed representations of $I^{s s}$. Below, we use the convention that $s_{1}$ and $t_{1}$ means a vertex on row 1 (i.e. has coordinate (1,?)), and $s_{2}$ and $t_{2}$ means a vertex on row 2 (i.e. has coordinate ( $2, ?$ )).

Then $I$ is in one of the following four cases, and the value of the compressed multiplicity $\underline{d}_{M}^{s s}(I)=d_{M^{\prime}}\left(I^{\prime}\right)$ is given by the respective formula.

- If I is a rectangle with source $s$ and sink $t$ then

$$
d_{M^{\prime}}\left(I^{\prime}\right)=\operatorname{rank} M(s \rightarrow t)
$$

- If I has sources $s_{1}, s_{2}$ and sink $t_{2}$ then

$$
\begin{aligned}
d_{M^{\prime}}\left(I^{\prime}\right)= & \operatorname{rank} M\left(s_{2} \rightarrow t_{2}\right)+\operatorname{rank} M\left(s_{1} \rightarrow t_{2}\right) \\
& -\operatorname{rank}\left[M\left(s_{2} \rightarrow t_{2}\right) \quad M\left(s_{1} \rightarrow t_{2}\right)\right]
\end{aligned}
$$

- If I has source $s_{1}$ and sinks $t_{1}, t_{2}$ then

$$
d_{M^{\prime}}\left(I^{\prime}\right)=\operatorname{rank} M\left(s_{1} \rightarrow t_{2}\right)+\operatorname{rank} M\left(s_{1} \rightarrow t_{1}\right)-\operatorname{rank}\left[\begin{array}{l}
M\left(s_{1} \rightarrow t_{2}\right) \\
M\left(s_{1} \rightarrow t_{1}\right)
\end{array}\right]
$$

- If I has sources $s_{1}, s_{2}$ and sinks $t_{1}, t_{2}$ then

$$
\begin{aligned}
d_{M^{\prime}}\left(I^{\prime}\right)= & \operatorname{rank}\left[\begin{array}{cc}
M\left(s_{2} \rightarrow t_{2}\right) & M\left(s_{1} \rightarrow t_{2}\right) \\
0 & M\left(s_{1} \rightarrow t_{1}\right)
\end{array}\right]+\operatorname{rank} M\left(s_{1} \rightarrow t_{2}\right) \\
& -\operatorname{rank}\left[\begin{array}{l}
M\left(s_{1} \rightarrow t_{2}\right) \\
M\left(s_{1} \rightarrow t_{1}\right)
\end{array}\right]-\operatorname{rank}\left[M\left(s_{2} \rightarrow t_{2}\right) \quad M\left(s_{1} \rightarrow t_{2}\right)\right]
\end{aligned}
$$

Proof. Each element $I$ of $\mathbb{I}_{2, n}$ has a staircase form, which is denoted by:

$$
I=\bigsqcup_{i=j}^{k}\left[b_{i}, d_{i}\right]_{i}
$$

for some integers $1 \leq j \leq k \leq 2$ and some integers $1 \leq b_{i} \leq d_{i} \leq n$ for each $j \leq i \leq k$ such that $b_{i+1} \leq b_{i} \leq d_{i+1} \leq d_{i}$ for any $i \in\{j, \ldots, k-1\}$.

The two cases given by

- $j=k$, or
- $b_{1}=b_{2}$ and $d_{1}=d_{2}$
correspond to $I$ being a rectangle (with source $s=\left(j, b_{j}\right)$ and $\operatorname{sink} t=\left(j, d_{j}\right)$, or source $s=\left(1, b_{1}\right)$ and $\operatorname{sink}$ $t=\left(2, d_{2}\right)$, respectively). Here, Proposition 4.16 gives the formula for the compressed multiplicity.

Thus, we are left with the cases that $1=j<k=2$, and that $b_{1} \neq b_{2}$ or $d_{1} \neq d_{2}$. By the general restriction that $b_{2} \leq b_{1} \leq d_{2} \leq d_{1}$, we have the following three cases

- $b_{2}<b_{1} \leq d_{2}=d_{1}$. This corresponds to the case that $I$ has sources $s_{1}=\left(1, b_{1}\right), s_{2}=\left(2, b_{2}\right)$ and sink $t_{2}=\left(2, d_{2}\right)$, as illustrated below:

with $I^{\mathrm{ss}}: s_{1} \longrightarrow t_{2} \longleftarrow s_{2}$ emphasized. Then, the compressed representations are given by

$$
I^{\prime}: K \xrightarrow{1} K \stackrel{1}{\longleftarrow} K \text { and } M^{\prime}: M\left(s_{1}\right) \xrightarrow{M\left(s_{1} \rightarrow t_{2}\right)} M\left(t_{2}\right)^{M\left(s_{2} \rightarrow t_{2}\right)} M\left(s_{2}\right) .
$$

We note that $I^{\prime}$ is injective with socle given by

$$
\operatorname{soc} I^{\prime}: 0 \xrightarrow{0} K \stackrel{0}{\longleftrightarrow} 0 .
$$

Using [19, Theorem 3], we have

$$
d_{M^{\prime}}\left(I^{\prime}\right)=\operatorname{dim} \operatorname{Hom}\left(I^{\prime}, M^{\prime}\right)-\operatorname{dim} \operatorname{Hom}\left(I^{\prime} / \operatorname{soc} I^{\prime}, M^{\prime}\right) .
$$

A homomorphism $I^{\prime} \rightarrow M^{\prime}$ is given by triples $(x, y, z)$ such that

commutes. That is, $y=M\left(s_{2} \rightarrow t_{2}\right) z=M\left(s_{1} \rightarrow t_{2}\right) x$. In other words, the homomorphism space $\operatorname{Hom}\left(I^{\prime}, M^{\prime}\right)$ is given by solutions to

$$
\left[M\left(s_{2} \rightarrow t_{2}\right) \quad M\left(s_{1} \rightarrow t_{2}\right)\right]\left[\begin{array}{c}
z \\
-x
\end{array}\right]=0
$$

(with $y$ fully determined by $x$ ), which has dimension equal to

$$
\operatorname{dim} M\left(s_{2}\right)+\operatorname{dim} M\left(s_{1}\right)-\operatorname{rank}\left[M\left(s_{2} \rightarrow t_{2}\right) \quad M\left(s_{1} \rightarrow t_{2}\right)\right] .
$$

On the other hand, a homomorphism $I^{\prime} / \operatorname{soc} I^{\prime} \rightarrow M^{\prime}$ is given by triples $(x, 0, z)$ such that

commutes. Thus

$$
\begin{aligned}
\operatorname{dim} \operatorname{Hom}\left(I^{\prime} / \operatorname{soc} I^{\prime}, M^{\prime}\right)= & \left(\operatorname{dim} M\left(s_{1}\right)-\operatorname{rank} M\left(s_{1} \rightarrow t_{2}\right)\right) \\
& +\left(\operatorname{dim} M\left(s_{2}\right)-\operatorname{rank} M\left(s_{2} \rightarrow t_{2}\right)\right) .
\end{aligned}
$$

Combining the above formulas yields the claimed formula for $d_{M^{\prime}}\left(I^{\prime}\right)$.

- $b_{2}=b_{1} \leq d_{2}<d_{1}$. This corresponds to the case that $I$ has source $s_{1}=\left(1, b_{1}\right)$ and sinks $t_{1}=$ $\left(1, d_{1}\right), t_{2}=\left(2, d_{2}\right)$ as illustrated below:


The proof for the formula of $d_{M^{\prime}}\left(I^{\prime}\right)$ in this case is dual to the previous case.

- $b_{2}<b_{1} \leq d_{2}<d_{1}$. This corresponds to the case that $I$ has sources $s_{1}=\left(1, b_{1}\right), s_{2}=\left(2, b_{2}\right)$ and sinks $t_{1}=\left(1, d_{1}\right), t_{2}=\left(2, d_{2}\right)$ as illustrated below:


Then, the compressed representations are given by

$$
I^{\prime}: K \xrightarrow{1} K \stackrel{1}{\longleftrightarrow} K \xrightarrow{1} K
$$

and

$$
M^{\prime}: M\left(s_{2}\right) \xrightarrow{M\left(s_{2} \rightarrow t_{2}\right)} M\left(t_{2}\right) \stackrel{M\left(s_{1} \rightarrow t_{2}\right)}{\rightleftarrows} M\left(s_{1}\right) \xrightarrow{M\left(s_{1} \rightarrow t_{1}\right)} M\left(t_{1}\right) .
$$

The almost split sequence starting from $I^{\prime}$ is given by

$$
0 \longrightarrow I^{\prime} \longrightarrow B \longrightarrow C \longrightarrow 0
$$

where

$$
B: K \xrightarrow{1} K \stackrel{\left[\begin{array}{ll}
1 & 0
\end{array}\right]}{\stackrel{ }{\longleftrightarrow}} K^{2} \xrightarrow{\left[\begin{array}{ll}
0 & 1
\end{array}\right]} K
$$

and

$$
C: 0 \xrightarrow{0} 0 \stackrel{0}{\longleftrightarrow} K \xrightarrow{0} 0 .
$$

Using [19, Theorem 3], we have

$$
\begin{equation*}
d_{M^{\prime}}\left(I^{\prime}\right)=\operatorname{dim} \operatorname{Hom}\left(I^{\prime}, M^{\prime}\right)-\operatorname{dim} \operatorname{Hom}\left(B, M^{\prime}\right)+\operatorname{dim} \operatorname{Hom}\left(C, M^{\prime}\right) \tag{6.1}
\end{equation*}
$$

For $(x, y, z, w) \in \operatorname{Hom}\left(I^{\prime}, M^{\prime}\right)$, the commutativity of

is equivalent to

$$
M\left(s_{2} \rightarrow t_{2}\right) x=y=M\left(s_{1} \rightarrow t_{2}\right) z \text { and } w=M\left(s_{1} \rightarrow t_{1}\right) z
$$

Then, each homomorphism is uniquely determined by a solution of

$$
\left[M\left(s_{2} \rightarrow t_{2}\right) \quad M\left(s_{1} \rightarrow t_{2}\right)\right]\left[\begin{array}{c}
x \\
-z
\end{array}\right]=0
$$

Thus,

$$
\begin{align*}
\operatorname{dim} \operatorname{Hom}\left(I^{\prime}, M^{\prime}\right)= & \operatorname{dim} M\left(s_{2}\right)+\operatorname{dim} M\left(s_{1}\right) \\
& -\operatorname{rank}\left[M\left(s_{2} \rightarrow t_{2}\right) \quad M\left(s_{1} \rightarrow t_{2}\right)\right] . \tag{6.2}
\end{align*}
$$

Next, for $(x, y, z, w) \in \operatorname{Hom}\left(B, M^{\prime}\right)$, the commutativity of

is equivalent to

$$
\begin{aligned}
M\left(s_{2} \rightarrow t_{2}\right) x & =y \\
M\left(s_{1} \rightarrow t_{2}\right) z_{1} & =y \\
M\left(s_{1} \rightarrow t_{2}\right) z_{2} & =0 \\
M\left(s_{1} \rightarrow t_{1}\right) z_{1} & =0 \\
M\left(s_{1} \rightarrow t_{1}\right) z_{2} & =w .
\end{aligned}
$$

Rewriting the above, we get that each homomorphism is uniquely determined by a solution to

$$
\left[\begin{array}{cc}
M\left(s_{2} \rightarrow t_{2}\right) & M\left(s_{1} \rightarrow t_{2}\right) \\
0 & M\left(s_{1} \rightarrow t_{1}\right)
\end{array}\right]\left[\begin{array}{c}
x \\
-z_{1}
\end{array}\right]=0 \text { and } M\left(s_{1} \rightarrow t_{2}\right) z_{2}=0
$$

with $y$ and $w$ determined from $x$ and $z_{2}$, respectively. Thus,

$$
\begin{align*}
\operatorname{dim} \operatorname{Hom}\left(B, M^{\prime}\right)= & \operatorname{dim} M\left(s_{2}\right)+\operatorname{dim} M\left(s_{1}\right) \\
& -\operatorname{rank}\left[\begin{array}{cc}
M\left(s_{2} \rightarrow t_{2}\right) & M\left(s_{1} \rightarrow t_{2}\right) \\
0 & M\left(s_{1} \rightarrow t_{1}\right)
\end{array}\right]  \tag{6.3}\\
& +\operatorname{dim} M\left(s_{1}\right)-\operatorname{rank} M\left(s_{1} \rightarrow t_{2}\right)
\end{align*}
$$

Finally, it is clear that

$$
\begin{align*}
\operatorname{dim} \operatorname{Hom}\left(C, M^{\prime}\right) & =\operatorname{dim}\left(\operatorname{ker} M\left(s_{1} \rightarrow t_{1}\right) \cap \operatorname{ker} M\left(s_{1} \rightarrow t_{2}\right)\right) \\
& =\operatorname{dim} \operatorname{ker}\left[\begin{array}{l}
M\left(s_{1} \rightarrow t_{1}\right) \\
M\left(s_{1} \rightarrow t_{2}\right)
\end{array}\right]  \tag{6.4}\\
& =\operatorname{dim} M\left(s_{1}\right)-\operatorname{rank}\left[\begin{array}{l}
M\left(s_{1} \rightarrow t_{1}\right) \\
M\left(s_{1} \rightarrow t_{2}\right)
\end{array}\right] .
\end{align*}
$$

Substituting Equations (6.2), (6.3), (6.4) into Equation (6.1) gives the claimed formula.

Let $\omega<2.373$ be the matrix multiplication exponent [23, 24].
Theorem 6.5 (Compressed multiplicity $\left(2 \times n\right.$ case)). For $M$ a persistence module over $\vec{G}_{2, n}$, Algorithm 3 computes $\underline{d}_{M}^{s s}$ with time complexity

$$
O\left(\frac{2^{\omega}+5}{24} n^{4} d^{\omega}\right)
$$

where $d=\max _{v \in\left(\vec{G}_{2, n}\right)_{0}} \operatorname{dim} M(v)$.

Proof. First, let us analyze Line 3 of Algorithm 3, which computes $M((x, y) \rightarrow(i, j))=M(p)$ for $(x, y) \neq$ $(i, j)$ with a path $p$ from $(x, y)$ to $(i, j)$. The value of $M(p)$ for paths $p$ with length equal to 1 (arrows) are already known. Assume that the values of $M(p)$ for all paths of length $\ell$ are already computed. Then, the value of $M(p)$ for each path $p$ of length $\ell+1$ can be computed by one matrix multiplication each. We note further that $M((x, y) \rightarrow(i, j))=M(p)$ does not depend on the which particular path $p$ is taken from $(x, y)$ to $(i, j)$. Thus, we can inductively compute the value of $M((x, y) \rightarrow(i, j))$ using one matrix multiplication for each pair of vertices $(x, y),(i, j)$ such that $(x, y) \neq(i, j)$ and there is a path of length greater than 1 from $(x, y)$ to $(i, j)$. Since there are

$$
\frac{3}{2}(n+1) n-2 n-(3 n-2)=O\left(\frac{3}{2} n^{2}\right)
$$

such pairs of vertices $(x, y) \neq(i, j)$ in the $2 \times n$ commutative grid by a simple combinatorial argument, Line 3 of Algorithm 3 can be performed in $O\left(\frac{3}{2} n^{2} d^{\omega}\right)$.

Next, we analyze Lines 4 to 7 of Algorithm 3. By [8, Corollary 4.12], there are

$$
\# \mathbb{I}_{2, n}=\frac{1}{24} n(n+1)\left(n^{2}+5 n+30\right)=O\left(\frac{1}{24} n^{4}\right)
$$

intervals $I$ to process. For each interval $I$, the computation of $d_{M^{\prime}}\left(I^{\prime}\right)$ using Proposition 6.4 involves computing the rank of a $2 d \times 2 d$, a $2 d \times d$, a $d \times 2 d$, and a $d \times d$ matrix in the worst case. Note that the rank of an $e \times f$ matrix $(e \leq f)$ can be computed with $O\left(f e^{\omega-1}\right)$ field operations by Gaussian elimination [25]. Thus, we get a cost of $O\left(\left(2^{\omega} d^{\omega}+5 d^{\omega}\right) \frac{1}{24} n^{4}\right)$ for the computation of $d_{M^{\prime}}\left(I^{\prime}\right)$.

Overall, we get a cost of $O\left(\frac{3}{2} n^{2} d^{\omega}+\frac{2^{\omega}+5}{24} n^{4} d^{\omega}\right)$ dominated by the latter term, giving the result.
For $I \in \mathbb{I}_{2, n}$, as shown in Example 3.3, $\# \operatorname{Cov}(I) \leq 4$. Thus, we get the following.
Corollary 6.6 (Möbius inversion $\delta_{M}^{*}(2 \times n$ case $)$ ). With $m=2$, Algorithm 2 (Möbius inversion $\delta_{M}^{*}$ of $\underline{d}_{M}^{*}$ ) can be performed with time complexity

$$
O\left(\frac{16}{3} n^{4}\right)
$$

Proof. Substituting $m=2, C=4$, and $\# \mathbb{I}_{2, n}=O\left(\frac{1}{24} n^{4}\right)$ into

$$
O\left(\# \mathbb{I}_{m, n} 2^{C} C \min \{m, n\}\right)
$$

from Theorem 6.3, we get the result.
Combining Theorem 6.5 and Corollary 6.6 with $*=$ ss, we get an overall cost of

$$
O\left(\frac{2^{\omega}+5}{24} n^{4} d^{\omega}+\frac{16}{3} n^{4}\right)
$$

for computing the interval-decomposable approximation $\delta^{\mathrm{ss}}(M)$ of $M$ in the $2 \times n$ case.
Implementation. As part of the software "pmgap" [26], we provide an implementation of Algorithms 3 and 2 in the $2 \times n$ case. The software "pmgap" builds upon the GAP [27] package QPA [28], which provides data structures and algorithms for computations on (quotients of) path algebras and their representations. The software "pmgap" uses those data structures to represent equioriented commutative grids and persistence modules over them, and implements the algorithms of this paper not in QPA.

Randomly generated persistence modules. For the computational experiments below, given values for $n$ and $d$ we randomly generate persistence modules $V$ (with $\mathbb{F}_{2}$ coefficients) over the commutative grid $\vec{G}_{2, n}$, such that all the vector spaces of $V$ have dimension $d$.

Below, whenever we say to randomly generate a $j \times k \mathbb{F}_{2}$-matrix $M$, we simply generate a matrix with entries independently and uniformly sampled from $\mathbb{F}_{2}$. If required, it is also possible to randomly choose a valid rank and then generate a random matrix with that rank. However, this comes at the cost of more computation time to generate the random matrices.

We use the following procedure to randomly generate the persistence module $V$. First, we randomly generate $d \times d \mathbb{F}_{2}$-matrices for each of the solid arrows below:


Then, for each square from right to left, we iteratively compute pullbacks (to guarantee commutativity) and multiply with another random matrix (to reach the correct dimension $d$ and to add more randomness). That is, given $d \times d$ matrices representing the linear maps $f$ and $g$ as below:

we compute (matrices with respect to some some basis of) the pullback maps ( $\phi_{1}, \phi_{2}$ ). Then, we randomly generate a $k \times d$ matrix representing $\phi_{3}$, and obtain the commutative diagram


It is clear that a persistence module $V$ over $\vec{G}_{2, n}$ is obtained by the above.
Computational experiments. We measure the time needed to compute the interval approximation using pmgap for some small values of $n$ and $d$. Computations were performed on Ubuntu 20.04.2 LTS running in WSL1 inside a Windows 10 Pro machine with an AMD Ryzen 5 5600X 6-Core ${ }^{5}$ Processor. In Table 2, we display the resulting runtimes in milliseconds. Each timing (each entry in the table) is measured as the average of at least five runs. Each run consists of the computation of the compressed multiplicity and interval approximation of a given persistence module. Additional runs are performed as needed so that the total time taken exceeds 100 ms , to ensure reliable measurement of runtimes; this is only needed for the smaller values of $n$ and $d$. Note that we exclude the time taken for generating the underlying path algebra, list of interval representations, and the persistence modules $V$.

For completeness, we also time the following operations: generation of the underlying path algebra and its list of interval representations, generation of a random persistence module, computation of the interval approximation. The results (of just one run each) are displayed in Table 3. Note that the underlying path algebra and its list of interval representations do not depend on the dimension $d$. Thus, we time that operation only once for each $n$.

[^4]Table 2: Runtimes (in ms) for the interval approximation using pmgap

| ${ }_{n}^{d}$ | 100 | 200 | 400 | 800 |
| :---: | :---: | :---: | :---: | :---: |
| 4 | 11.88 | 34.40 | 112.40 | 471.80 |
| 8 | 131.20 | 328.20 | 1,152.80 | 5,115.60 |
| 16 | 1,881.40 | 4,415.80 | 14,918.80 | 67,171.60 |

* Runtimes are measured as an average of at least five runs.
* Runtimes do not include time needed for generating the underlying path algebra, list of interval representations, and the persistence modules.

Table 3: Runtimes (in ms) in pmgap

| operation | $\boldsymbol{n}$ |  |  |  |  |
| :---: | ---: | ---: | ---: | ---: | ---: |
| algebra | 4 |  |  |  | 31.0 |
| and | 8 |  |  |  | 562.0 |
| its intervals | 16 |  |  |  | $16,578.0$ |
|  |  |  |  |  |  |
|  | $\boldsymbol{d}$ | 100.0 | 200.0 | 400.0 | 800.0 |
| operation | $\boldsymbol{n}$ |  |  |  |  |
| random | 4 | 46.0 | 171.0 | 640.0 | 2594.0 |
| persistence | 8 | 94.0 | 375.0 | $1,422.0$ | $5,516.0$ |
| module | 16 | 219.0 | 781.0 | $3,047.0$ | $12,609.0$ |
| interval | 4 | 15.0 | 31.0 | 109.0 | 484.0 |
| approximation | 8 | 125.0 | 328.0 | $1,156.0$ | $5,141.0$ |
|  | 16 | $1,828.0$ | $4,422.0$ | $14,953.0$ | $67,218.0$ |

Demonstrations. The pmgap repository [26] contains demonstrations for these computations.
We also provide a browser-based implementation [29] demonstrating the computation of interval approximation of randomly generated persistence modules. Note that the browser-based demo [29] was developed separately of pmgap, and does not rely on the installation of pmgap.

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[^1]:    ${ }^{1}$ A persistence module is said to be thin if all of its vector spaces have dimension at most 1 . For example, interval persistence modules are thin.
    ${ }^{2}$ Note that there is a difference between our convention and theirs in the order of arrows in paths. Namely, the path $\alpha_{n} \cdots \alpha_{1}$ in this paper is written as $\alpha_{1} \cdots \alpha_{n}$ in their book

[^2]:    ${ }^{3}$ More precisely, we use the inclusion-exclusion principle for finite measures, where we note that $\left(U(I), 2^{U(I)}, f\right)$ is a finite measure space.

[^3]:    ${ }^{4}$ A non-split short exact sequence $(E): 0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0$ is called an almost split sequence starting at $X$ if both $X$ and $Z$ are indecomposable, and if for any homomorphism $h: X \rightarrow V$, either $h$ is a split monomorphism or the pushout of ( $E$ ) along $h$ splits.

[^4]:    ${ }^{5}$ Note that the current implementation does not take advantage of multiple cores or threads.

